

# Internet Technology

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Twitter: @drchuck



<https://www.coursera.org/course/insidetheinternet>

coursera

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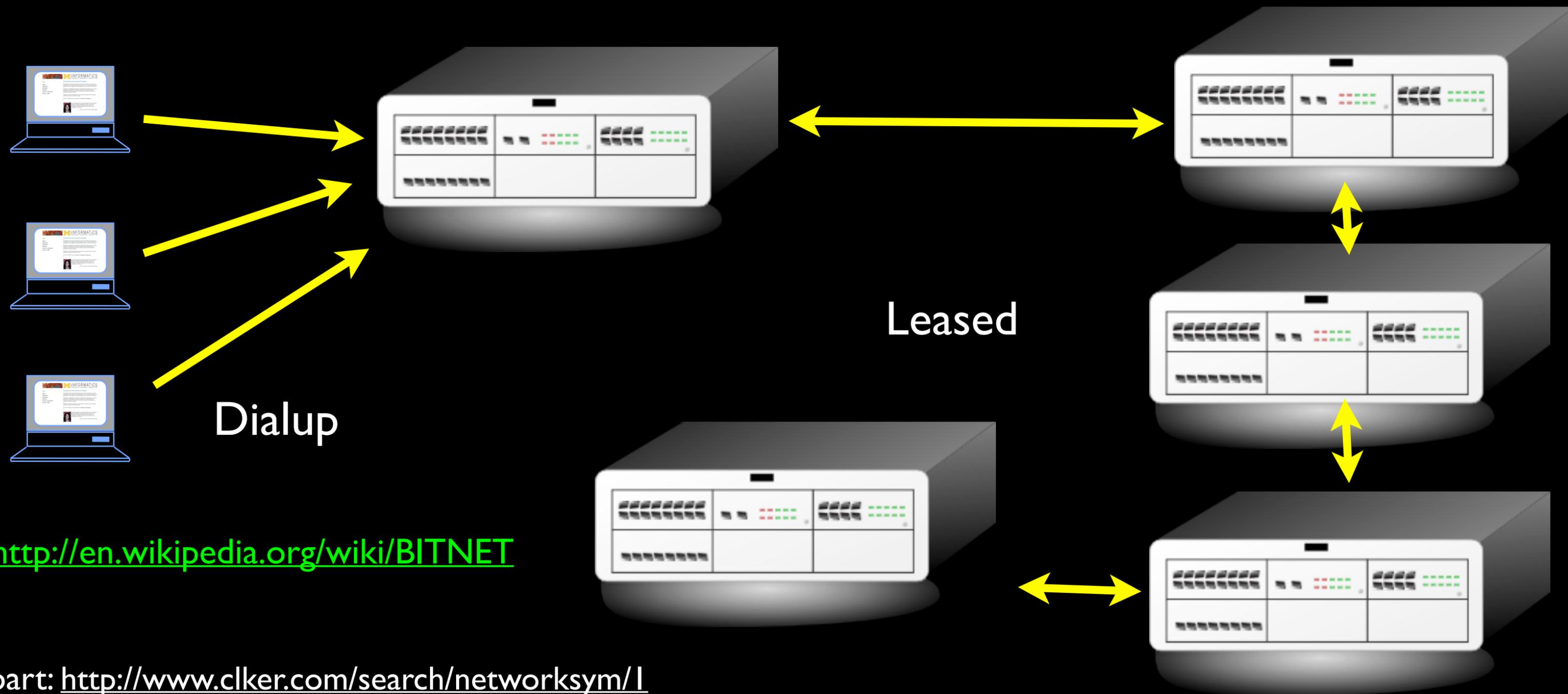
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BUT WHEN SHE TRACED THE  
KILLER'S IP ADDRESS... IT WAS  
IN THE 192.168/16 BLOCK!



<http://xkcd.com/742/>

# Store and Forward Networking

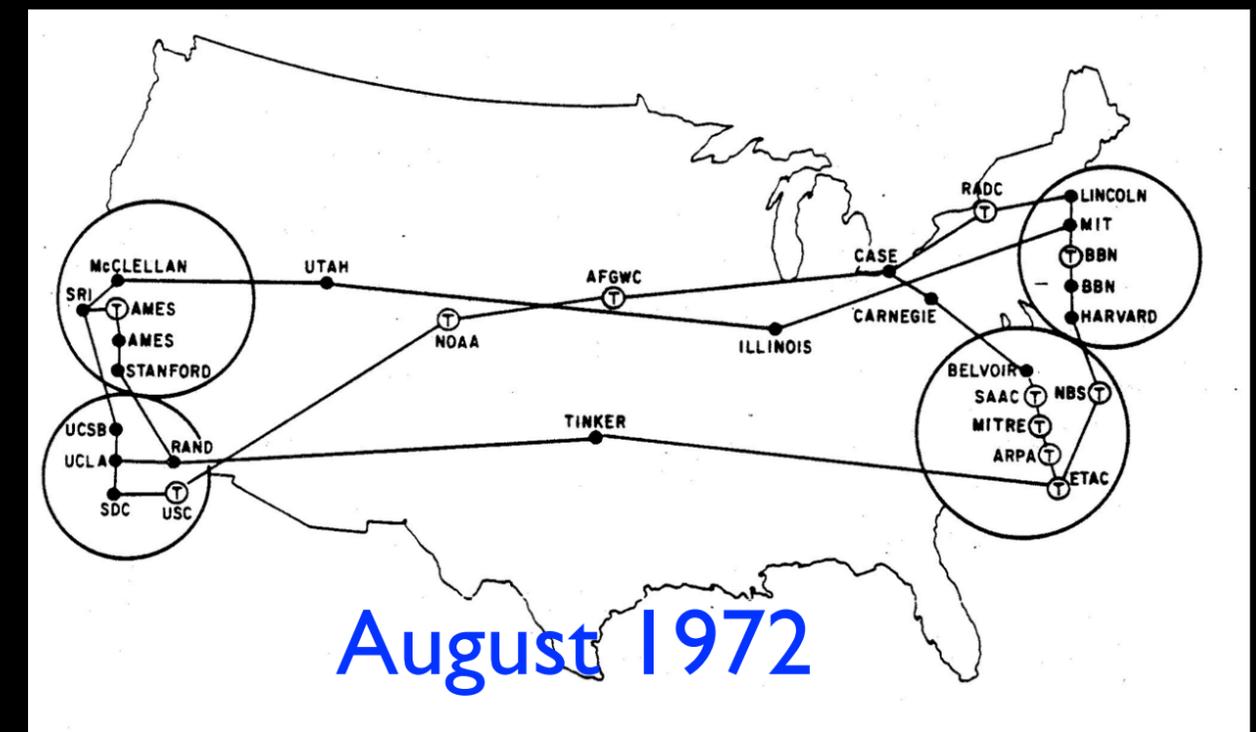
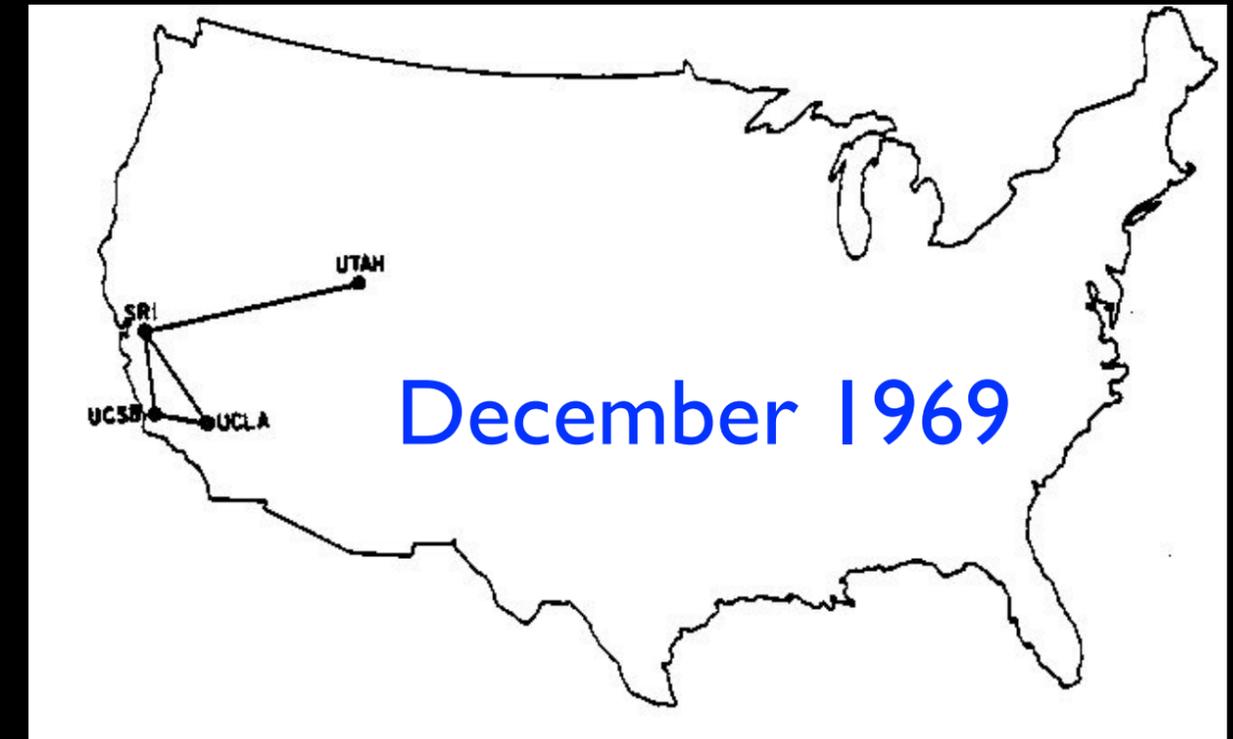


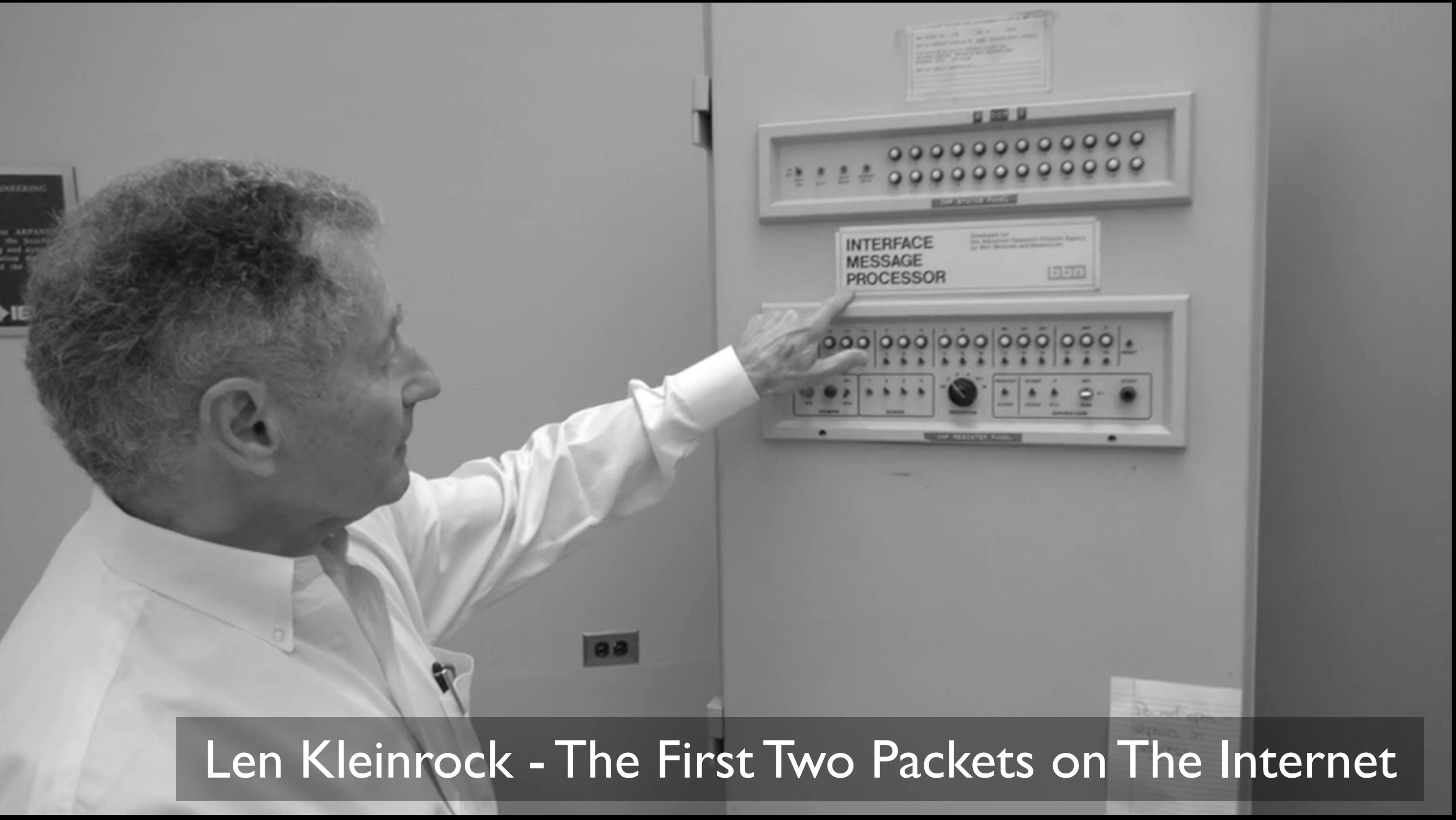
<http://en.wikipedia.org/wiki/BITNET>

Clipart: <http://www.clker.com/search/networksym/>

# Research Networks 1960-1980's

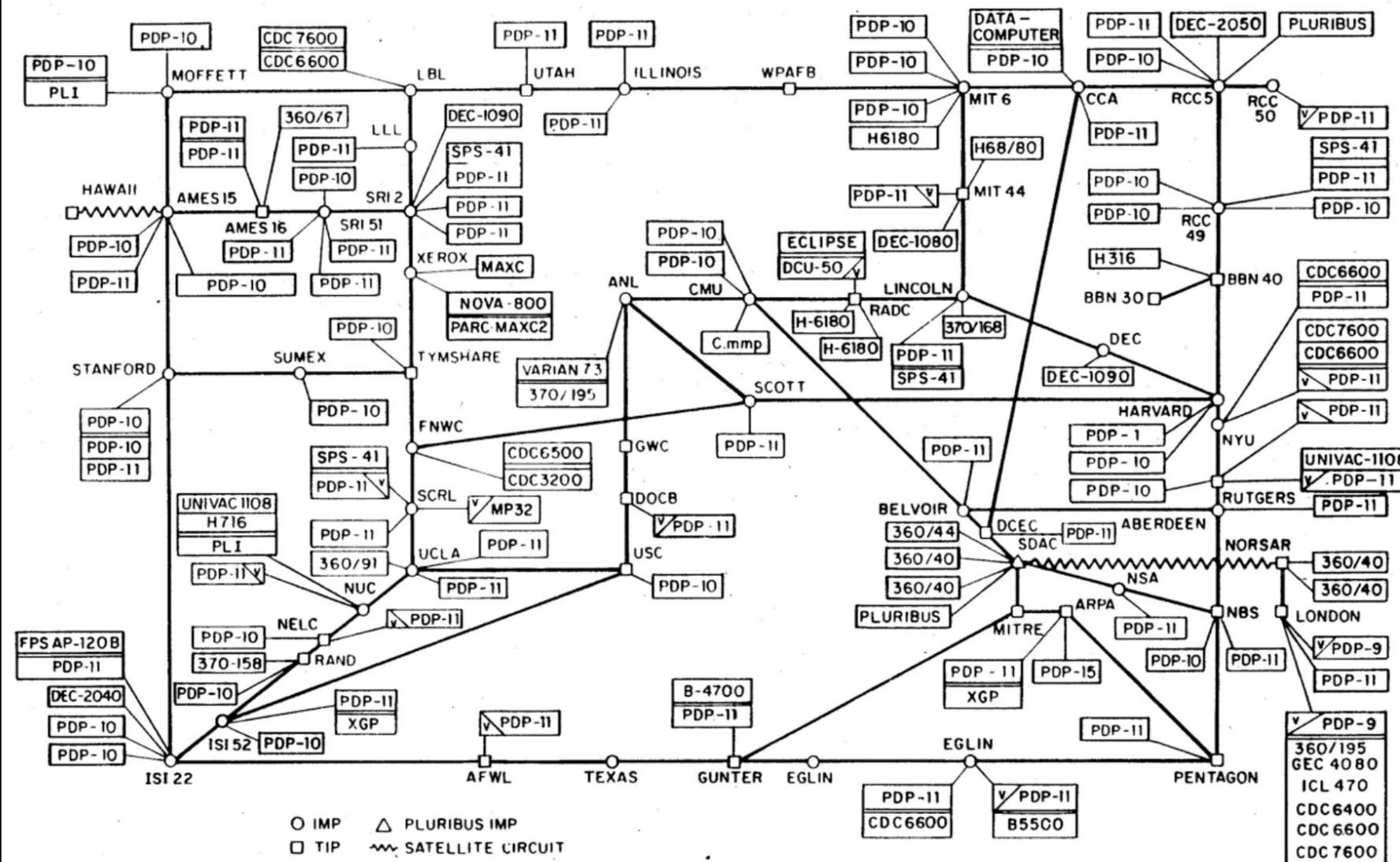
- How can we avoid having a direct connection between all pairs of computers?
- How to transport messages efficiently?
- How can we dynamically handle outages?





Len Kleinrock - The First Two Packets on The Internet

ARPANET LOGICAL MAP, MARCH 1977



(PLEASE NOTE THAT WHILE THIS MAP SHOWS THE HOST POPULATION OF THE NETWORK ACCORDING TO THE BEST INFORMATION OBTAINABLE, NO CLAIM CAN BE MADE FOR ITS ACCURACY)

NAMES SHOWN ARE IMP NAMES, NOT (NECESSARILY) HOST NAMES

Heart, F., McKenzie, A., McQuillian, J., and Walden, D., ARPANET Completion Report, Bolt, Beranek and Newman, Burlington, MA, January 4, 1978.

<http://som.csudh.edu/fac/lpress/history/arpamaps/arpametmar77.jpg>

# Efficient Message Transmission: Packet Switching

- Challenge: in a simple approach, like store-and-forward, large messages block small ones
- Break each message into **packets**
- Can allow the packets from a single message to travel over different paths, dynamically adjusting for use
- Use special-purpose computers, called **routers**, for the traffic control

# Packet Switching - Postcards

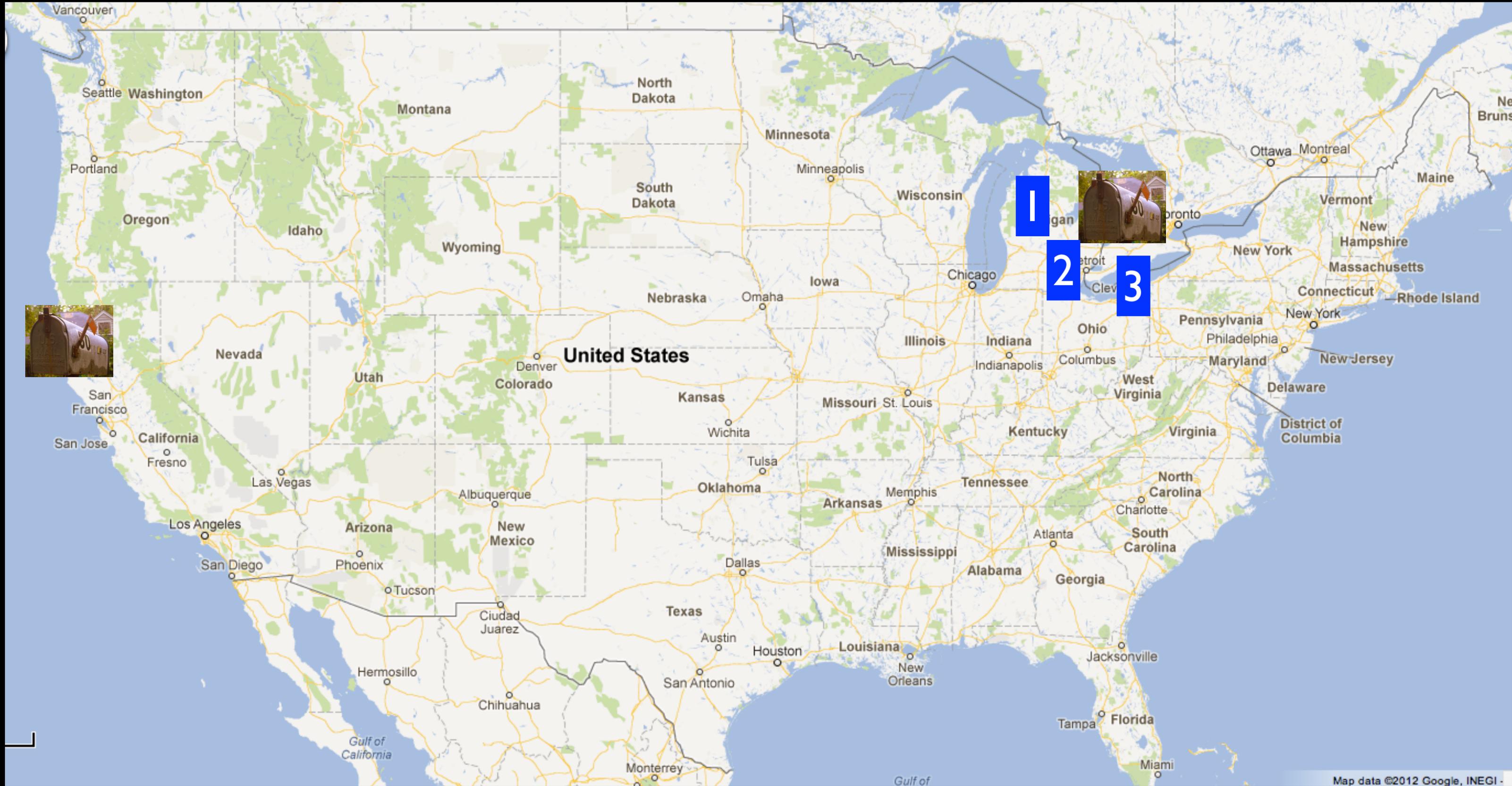
Hello there, have a nice day.

Hello ther (1, csev, daphne)

e, have a (2, csev, daphne)

nice day. (3, csev, daphne)



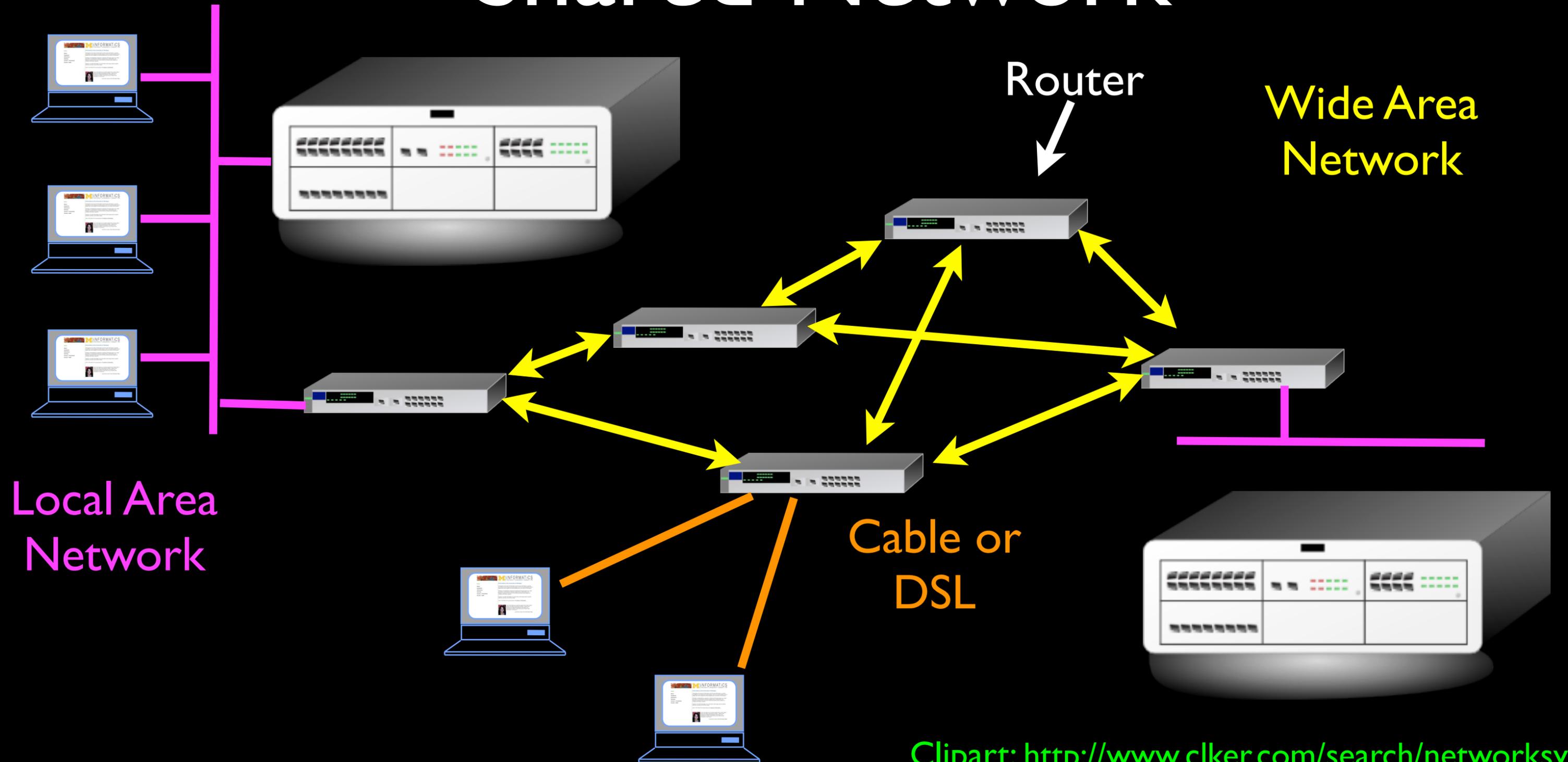


# Packet Switching - Postcards



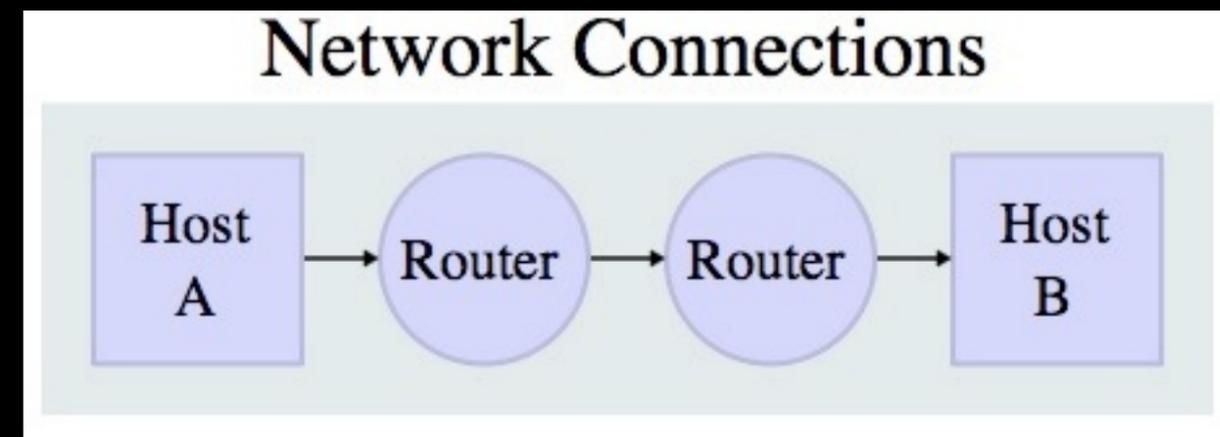
Hello there, have a nice day.

# Shared Network



# Shared Networks

- In order to keep cost low and the connections short geographically - data would be forwarded through several routers.
- Getting across the country usually takes about 10 “hops”
- Network designers continually add and remove links to “tune” their networks



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Layered Network Model

- A **layered** approach allows the problem of designing a network to be broken into more manageable sub problems
- Best-known model: **TCP/IP**—the “Internet Protocol Suite”
- There was also a 7 layer **OSI**: Open System Interconnection Model

Application Layer  
Web, E-Mail, File Transfer

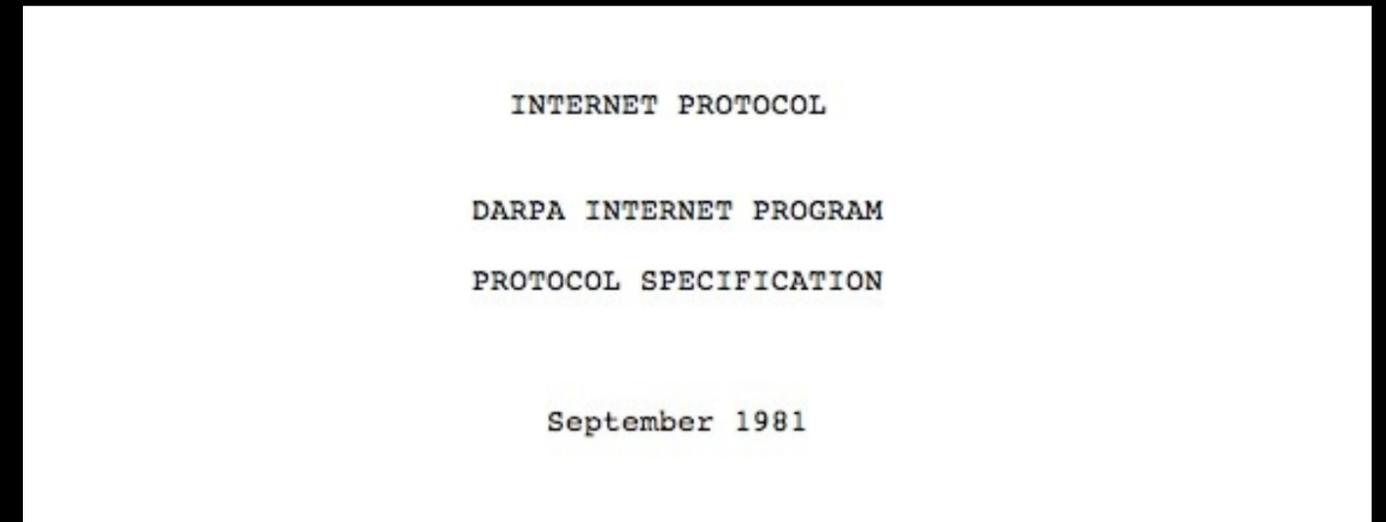
Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Internet Standards

- The standards for all of the Internet protocols (inner workings) are developed by an organization
- Internet Engineering Task Force (IETF)
- [www.ietf.org](http://www.ietf.org)
- Standards are called “RFCs” - “Request for Comments”



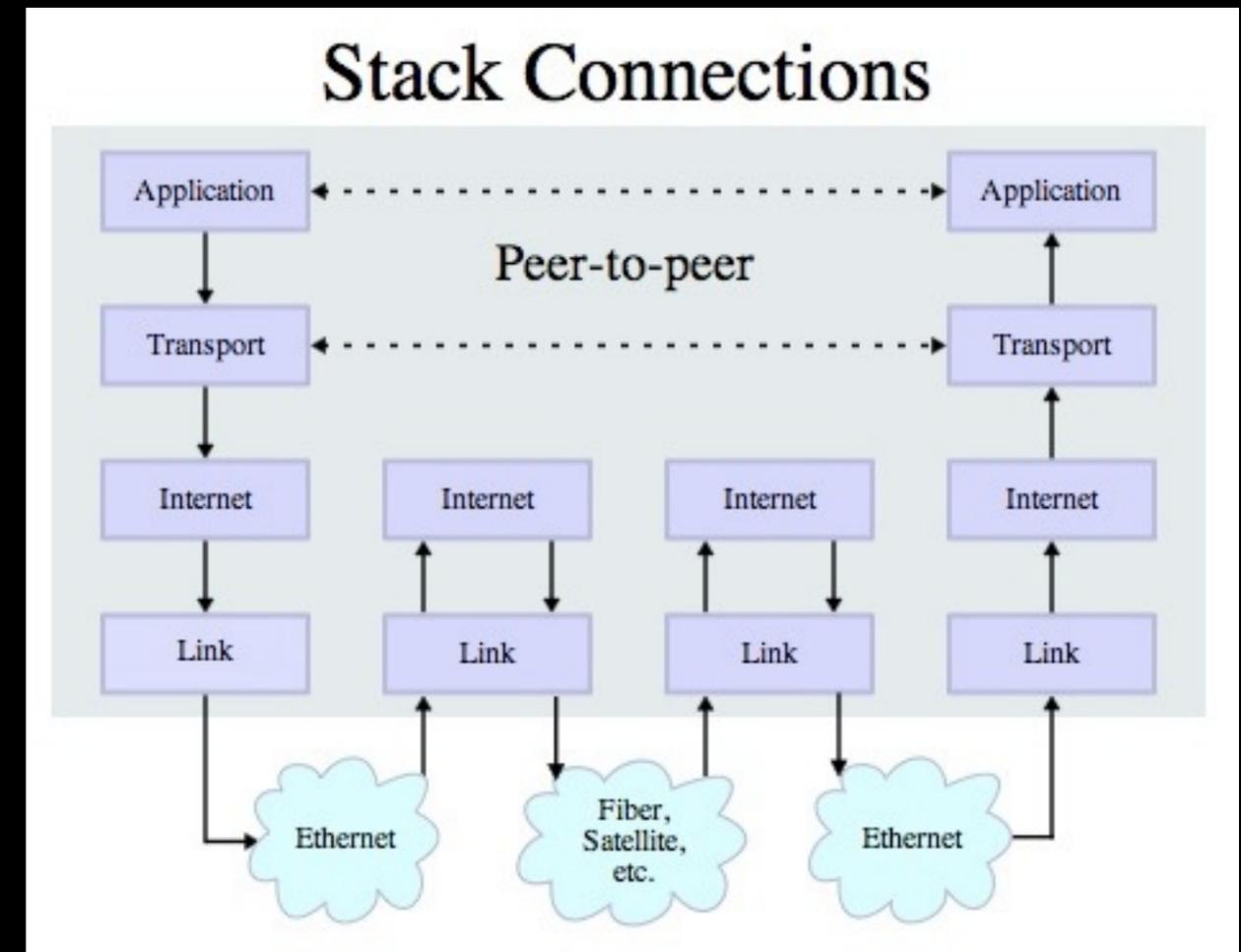
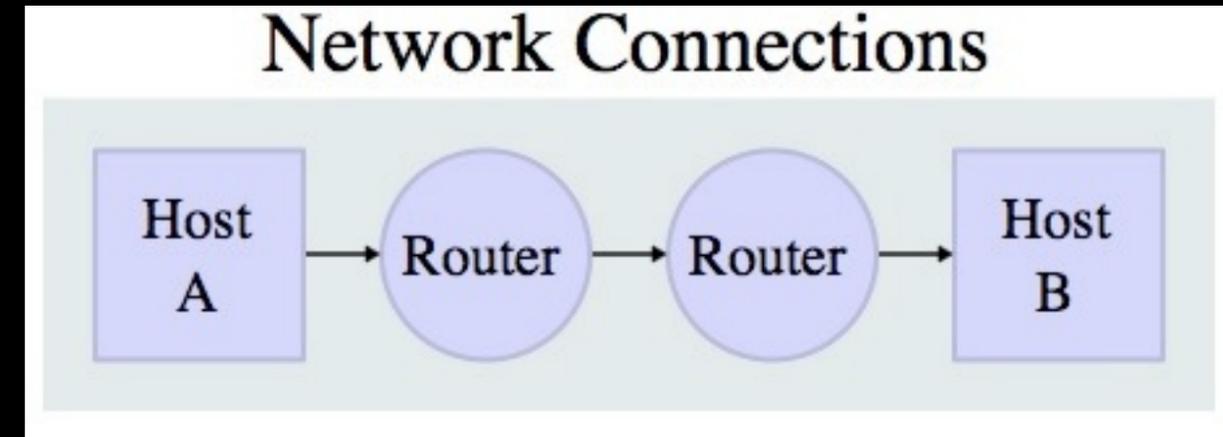
The internet protocol treats each internet datagram as an independent entity unrelated to any other internet datagram. There are no connections or logical circuits (virtual or otherwise).

The internet protocol uses four key mechanisms in providing its service: Type of Service, Time to Live, Options, and Header Checksum.

Source: <http://tools.ietf.org/html/rfc791>

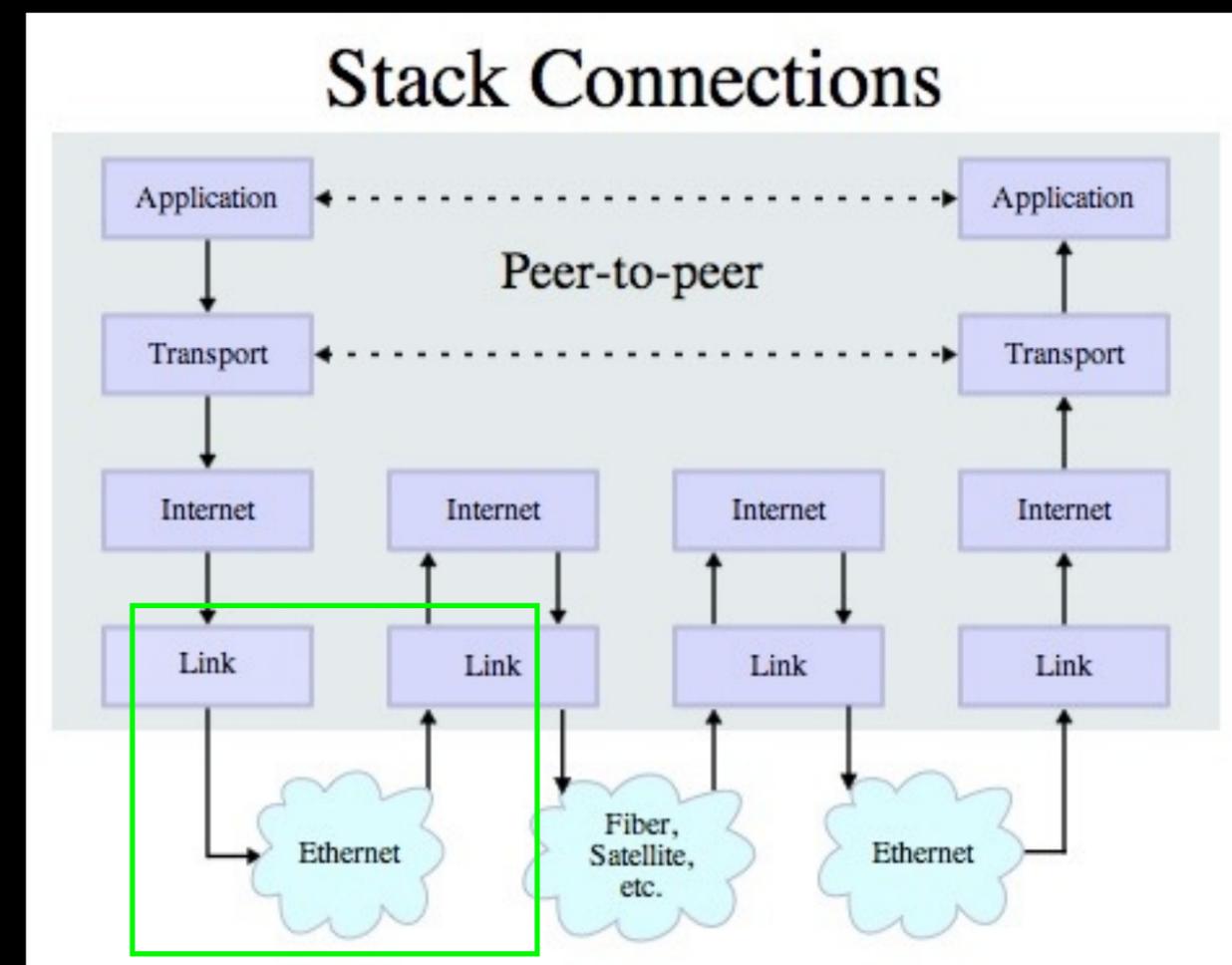
# Layered Architecture

- The Physical and Internet Layers are like trucks and trains - they haul stuff and get it to the right loading dock - it takes multiple steps
- The Transport layer checks to see if the trucks made it and send the stuff again if necessary



# Link Layer (aka Physical Layer)

- As your data crosses the country may use a different physical medium for each “hop”
- Wire, Wireless, Fiber Optic, etc.
- The link is “one hop” - Is it up or down? Connected or not?
- Very narrow focus - no view at all of the “whole Internet”

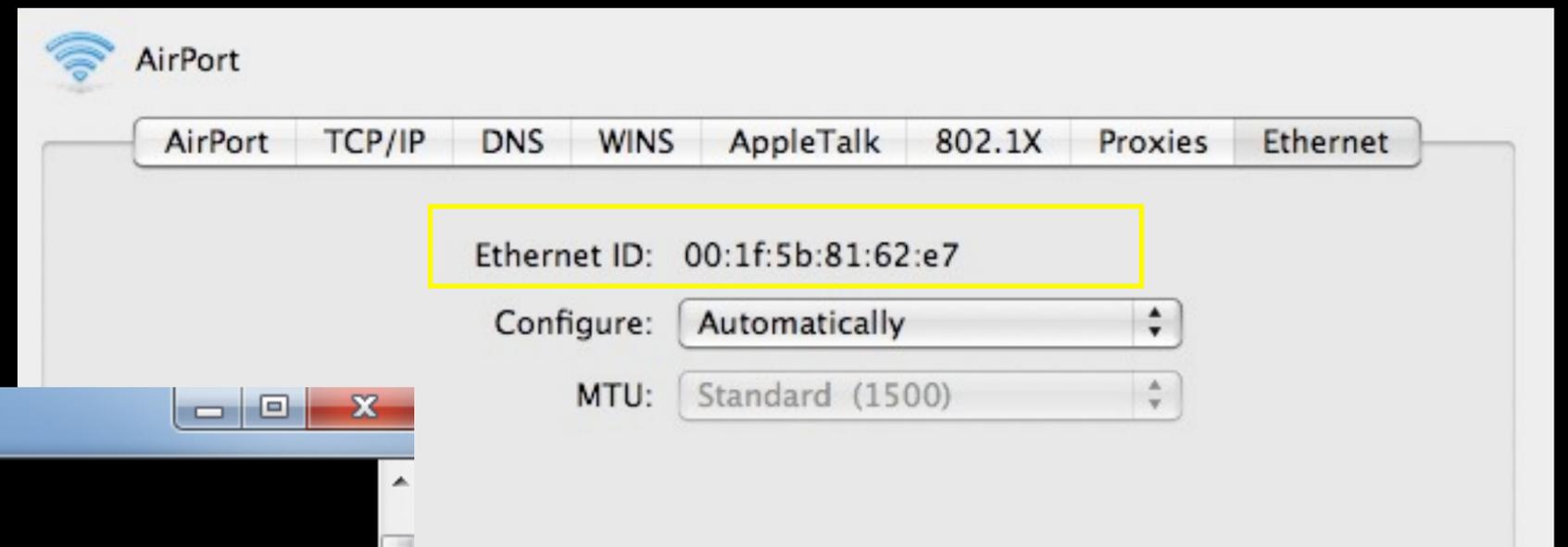


# Problems solved by the Link Layer

- **How does data get pushed onto a link?**
- **How is the link shared?**
- **Common Link Technologies**
  - **Ethernet**
  - **WiFi**
  - **Cable modem**
  - **DSL**
  - **Satellite**
  - **Optical**

# Link Layer Addresses

- Many physical layer devices have addresses built in to them by the manufacturer
  - Ethernet
  - Wireless Ethernet (Wifi)



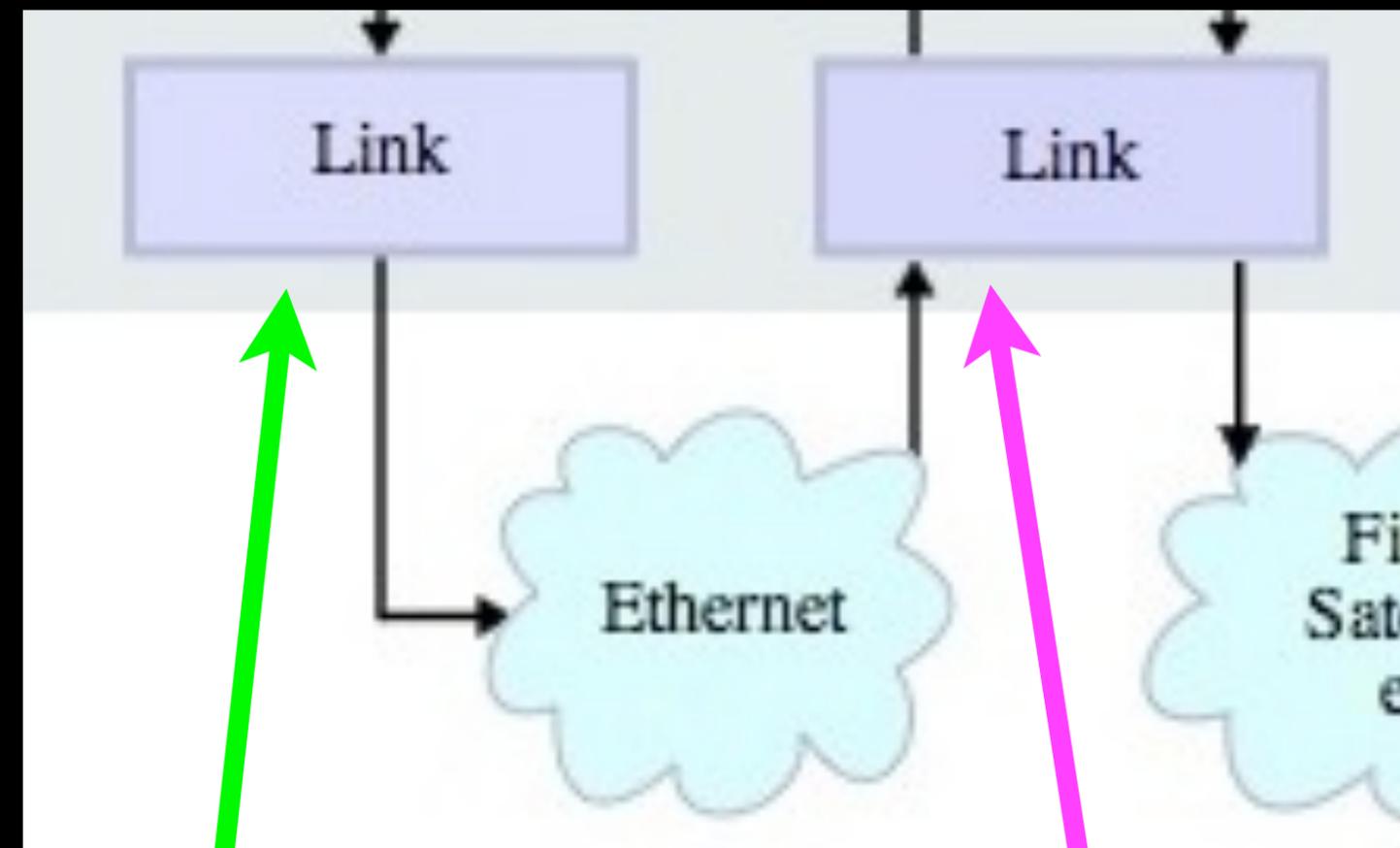
```
C:\Windows\system32\cmd.exe

Ethernet adapter Local Area Connection 2:

Connection-specific DNS Suffix . : umich.edu
Description . . . . . : Cisco Systems UPN Adapter
Physical Address. . . . . : 00-05-9A-3C-78-00
DHCP Enabled. . . . . : No
Autoconfiguration Enabled . . . . : Yes
```

# Link Layer

- Physical addresses are to allow systems to identify themselves on the ends of a single link
- Physical addresses go no farther than one link
- Sometimes links like Wifi and Wired Ethernet are shared with multiple computers



0f:21:63:12:b3:1a

98:2f:4e:78:01:b4

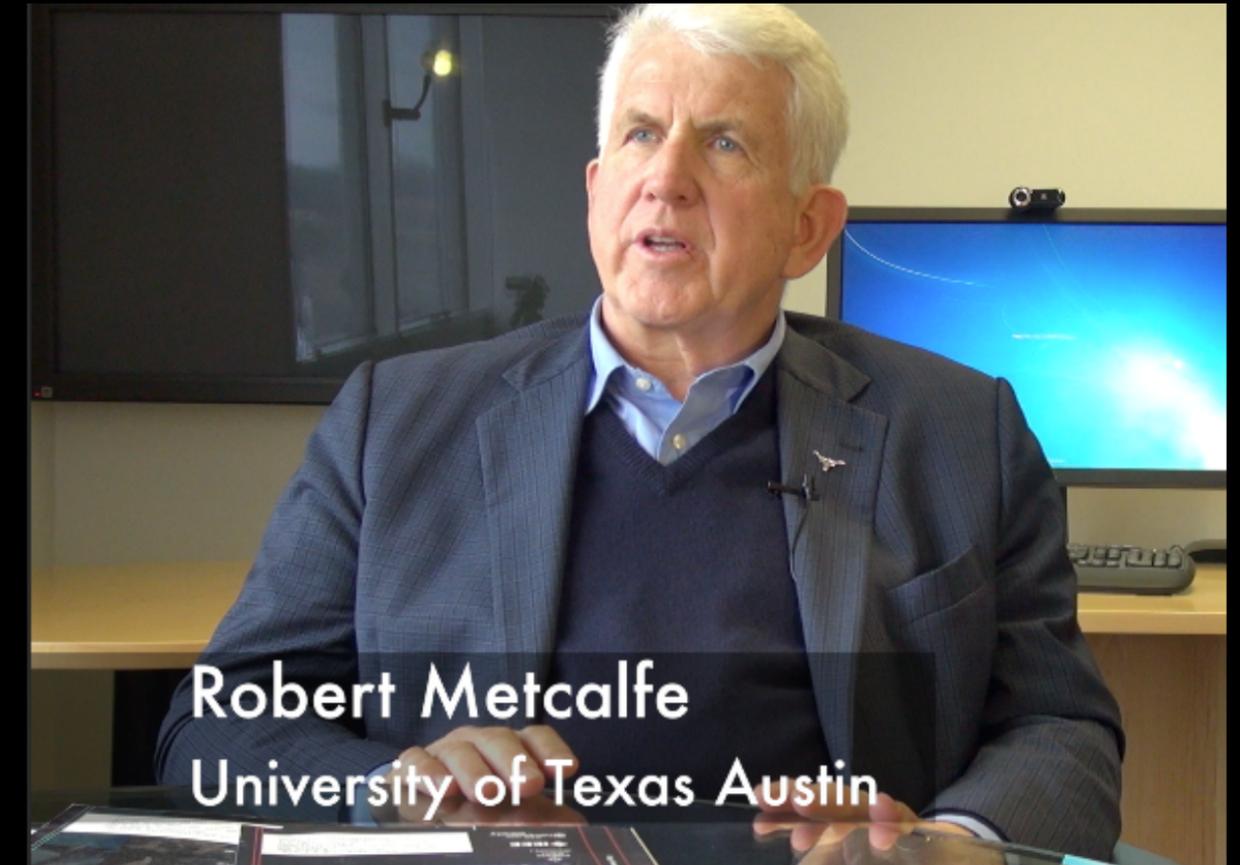
# Sharing Nicely - Avoiding Chaos

- **CSMA/CD Carrier Sense Multiple Access with Collision Detection**
- **To avoid garbled messages, systems must observe “rules” (Protocols)**
- **Ethernet rules are simple**
  - Wait for silence
  - Begin Transmitting data
  - Listen for your own data
  - If you cannot hear your own data clearly, assume a collision, stop and wait before trying again
  - Each system waits a different amount of time to avoid “too much politeness”



# Ethernet

- Invented at PARC (Xerox)
- The first Local-Area-Network
- Connected PC's to laser printers
- Inspired by an earlier wireless network called Aloha from the University of Hawaii



# Internetwork Layer (IP)

[http://en.wikipedia.org/wiki/Internet\\_Protocol](http://en.wikipedia.org/wiki/Internet_Protocol)

<http://en.wikipedia.org/wiki/Traceroute>

<http://en.wikipedia.org/wiki/Ping>

Application Layer  
Web, E-Mail, File Transfer

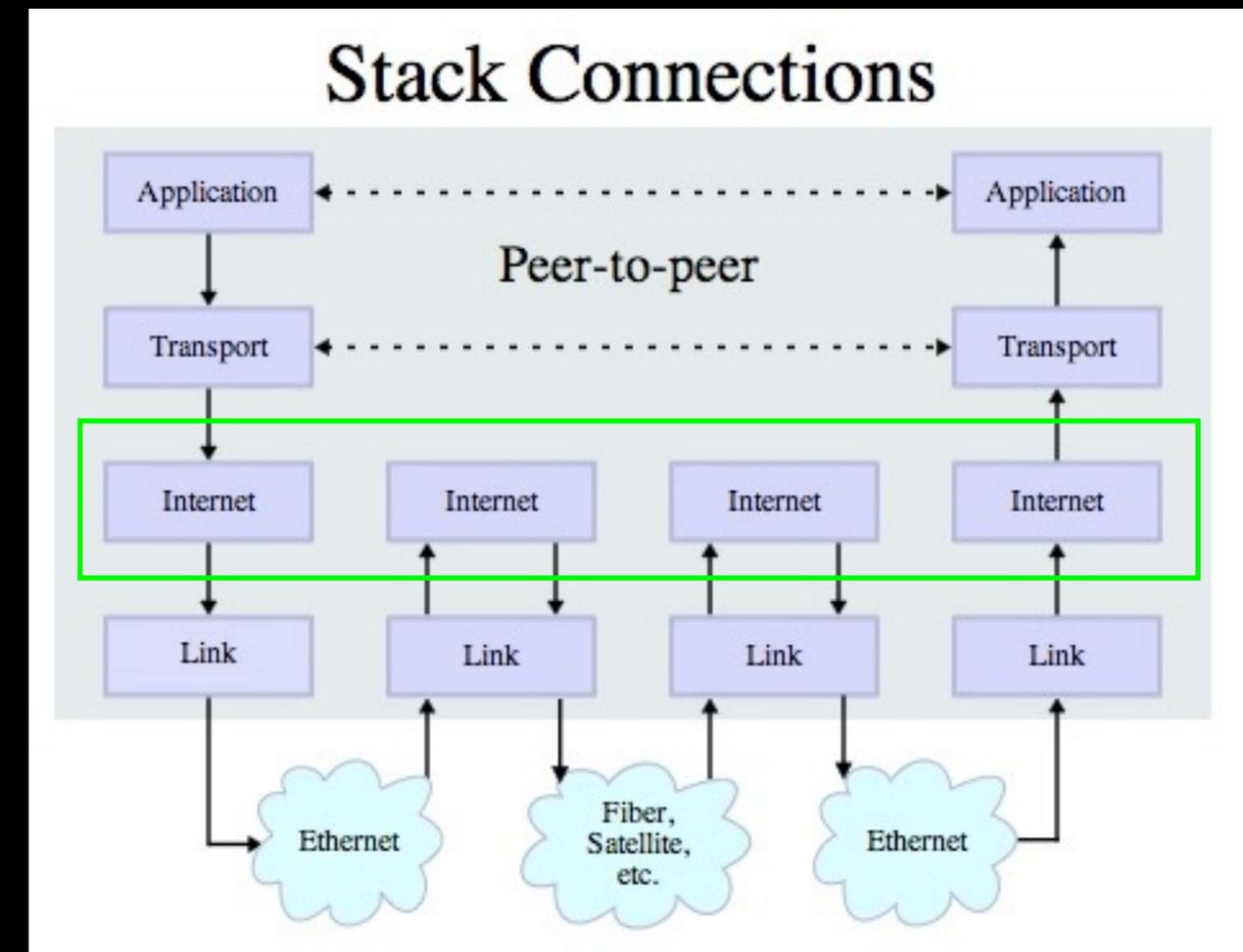
Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Internet Protocol Layer

- Goal: Gets your data from this computer to the other computer half way across the world
- Each router knows about nearby routers
- IP Is best effort - it is OK to drop data if things go bad...



# IP Addresses

- The IP address is the worldwide number which is associated with one particular workstation or server
- Every system which will send packets directly out across the Internet must have a unique IP address
- IP addresses are based on where station is connected
- IP addresses are not controlled by a single organization - address ranges are assigned
- They are like phone numbers – they get reorganized once in a great while

Google Custom Search

Search

Ads by Google

**Comcast Business Internet**

Get Comcast Business Class Today Find Local Offer For Your Business  
[www.ComcastBusinessService.com](http://www.ComcastBusinessService.com)

**IP Scan**

Scan your Network IP Free Network IP Scan.  
[IP.Scan.Qualys.com](http://IP.Scan.Qualys.com)

**Ip Address Providers**

Voice, Data, T-1, Mobile, Web Host. Get price quote. Special packages.  
[grow.cbeyond.net](http://grow.cbeyond.net)

**Reverse Email Lookup**

1) Type in Email Address. 2) Get Owner Name & Info.  
[EmailFinder.com](http://EmailFinder.com)

**Embarq is now CenturyLink**

Try One Voice & Data

What is my IP address?

BOOKMARK

Your IP address is **68.42.65.147**

(Now detects many proxy servers)

Free Trial.



- Tools
- [IP Lookup](#)
- [Blacklist Check](#)
- [Trace Email](#)
- [Visual Traceroute](#)
- [Traceroute](#)

68.42.65.147

Trace Now

IP Lookup now shows ISP, Organization, Proxy Status, and Connection Type!

IP Address Location: Ann Arbor, Michigan United States

Read about [GeoLocation accuracy](#).

What is an IP address?

Every device connected to the public Internet is assigned a unique number known as an Internet Protocol (IP) address. IP addresses consist of four numbers separated by periods (also called a 'dotted-quad') and look something like 127.0.0.1.

Since these numbers are usually assigned to internet service providers within region-based blocks, an IP address can often be used to identify the region or country from which a computer is connecting to the Internet. An IP address can sometimes be used to show the

# IP Address Format

- **Four numbers with dots - each number 1-255 (32 bits)**
- **Kind of like phone numbers with an “area code”**
- **The prefix of the address is “which network”**
- **While the data is traversing the Internet - all that matters is the network number**

(734) 764 1855  
Area code

Network  
Number

141.211.144.188  
141.211.\*.\*

While in the network, all that matters is the Network number.

141.211.144.188



To: 67.149.\*.\*



67.149.102.75

67.149.\*.\*

To: 67.149.94.33



67.149.94.33

To: 67.149.94.33

No single router knows the whole network - just which way to send data to get it "closer"

141.211.144.188



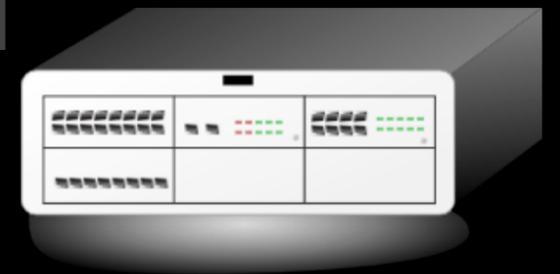
To: 67.149.\*.\*



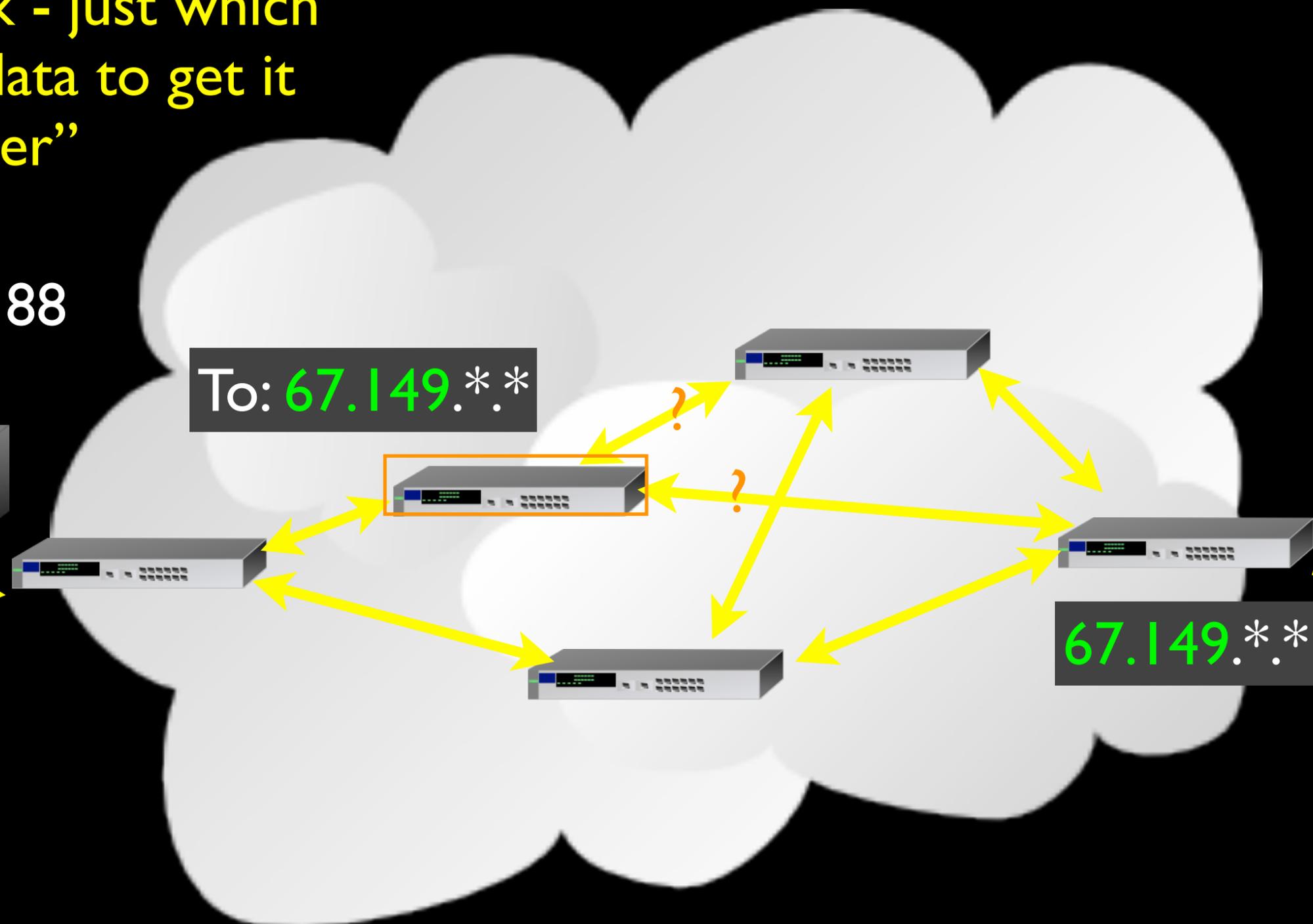
67.149.\*.\*



67.149.102.75



67.149.94.33

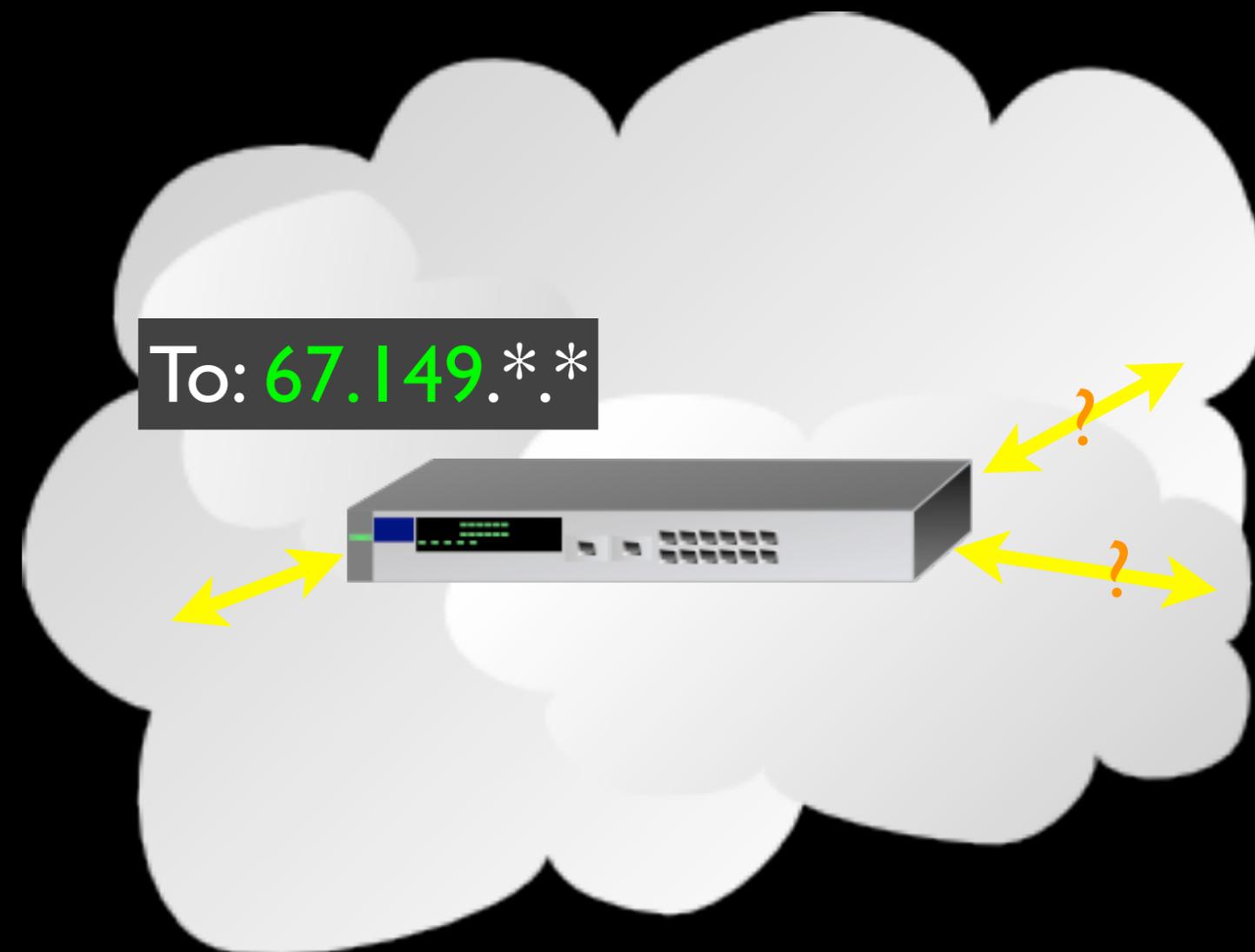


# Router Tables

Lists of where to send packets, based on destination network address;  
bandwidth on adjacent links;  
traffic on adjacent links;  
state of neighbor nodes (up or not);  
...

Updated dynamically

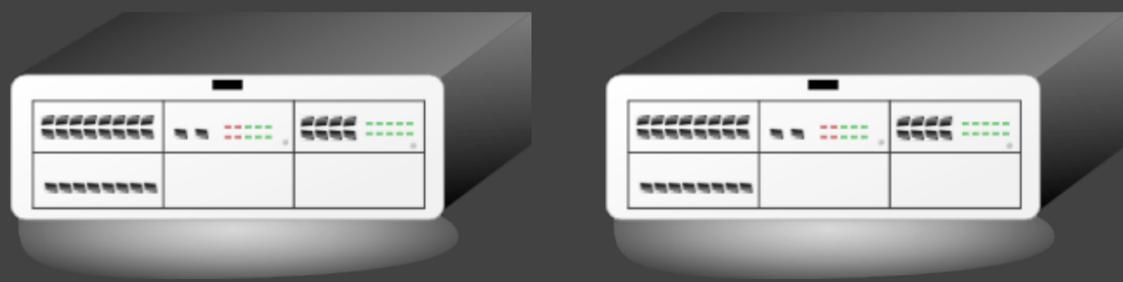
Routers “ask each other” for information



# IP Is Simple

Thousands of network connections.  
Billions of bytes of data per seconds.

Local Network



100's of servers

67.149.\*.\*

One "area code" to keep track of inside the Internet.



Thousands of user systems

# DHCP = Dynamic Host Configuration Protocol

Hello?

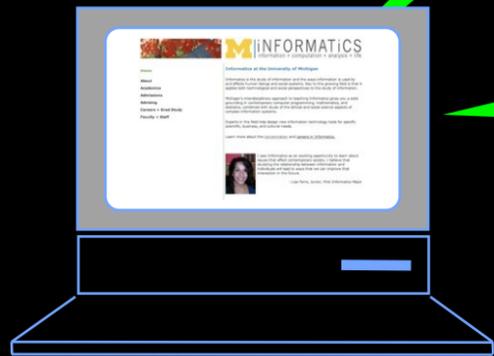
Here I am

What IP Address can I use?

141.26.14.1

141.26.14.1-100

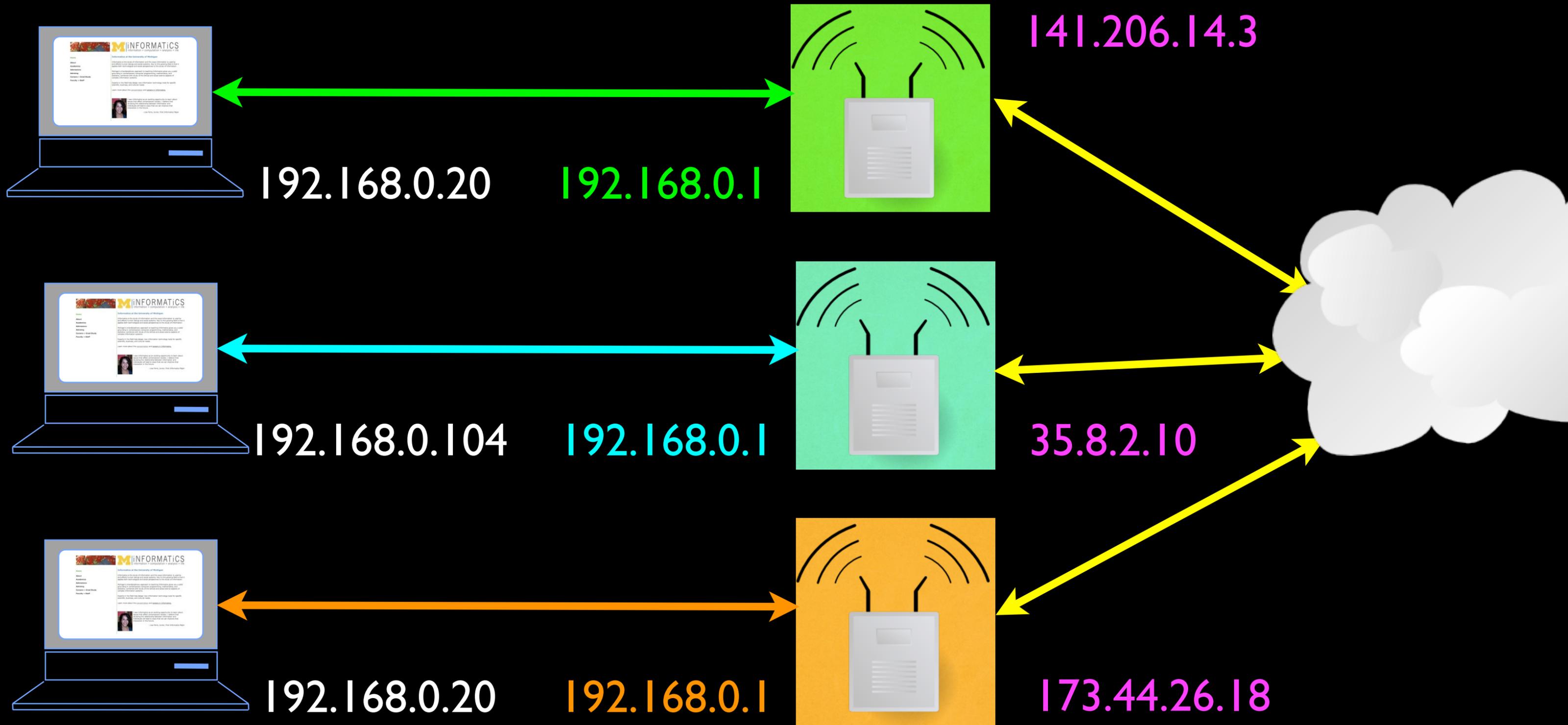
Use 141.26.14.7



# Non-Routable Addresses

- A typical home router does Network Address Translation (NAT)
- Your ISP gives your home router a real global routable address
- Your router gives out local addresses in a special range (192.168.\*.\*)
- The router maps remote addresses for each connection you make from within your home network

[http://en.wikipedia.org/wiki/Network\\_address\\_translation](http://en.wikipedia.org/wiki/Network_address_translation)



**NAT = Network Address Translation**

Clipart: <http://www.clker.com/search/networksym/>

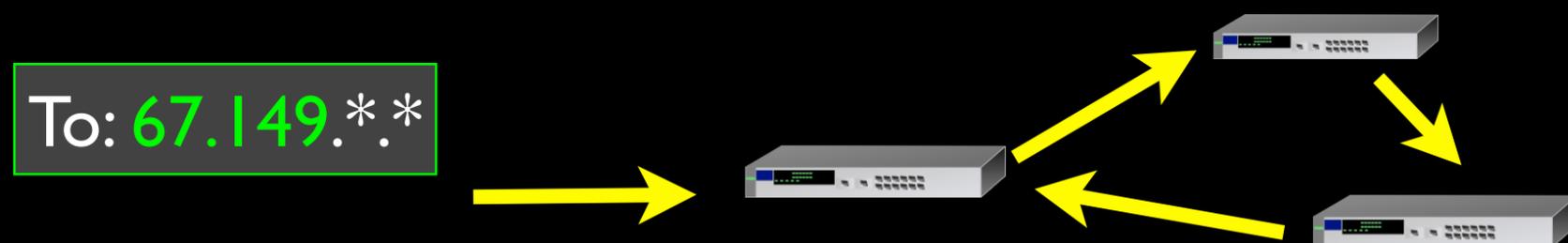
BUT WHEN SHE TRACED THE  
KILLER'S IP ADDRESS... IT WAS  
IN THE 192.168/16 BLOCK!



<http://xkcd.com/742/>

# Peering into the Internet

- Most systems have a command that will reveal the route taken across the internet (tracert on Windows and traceroute on Mac)
- Each IP packet has a field called “Time to Live” - TTL
- The TTL is used to deal with loops in the network - normally if routers got confused and ended up with a loop - the network would clog up rapidly.



# How Traceroute Works

- Normal packets are sent with a Time to Live (TTL) of 255 hops
- Trace route sends a packet with TTL=1, TTL=2, ...
- So each packet gets part-way there and then gets dropped and traceroute gets a notification of where the drop happens
- This builds a map of the nodes that a packet visits when crossing the Internet.

# Traceroute

```
$ traceroute www.stanford.edu
```

```
traceroute to www5.stanford.edu (171.67.20.37), 64 hops max, 40 byte packets
```

```
1 141.211.203.252 (141.211.203.252) 1.390 ms 0.534 ms 0.490 ms
2 v-bin-seb.r-bin-seb.umnet.umich.edu (192.122.183.61) 0.591 ms 0.558 ms 0.570 ms
3 v-bin-seb-i2-aa.merit-aa2.umnet.umich.edu (192.12.80.33) 6.610 ms 6.545 ms 6.654 ms
4 192.122.183.30 (192.122.183.30) 7.919 ms 7.209 ms 7.122 ms
5 so-4-3-0.0.rtr.kans.net.internet2.edu (64.57.28.36) 17.672 ms 17.836 ms 17.673 ms
6 so-0-1-0.0.rtr.hous.net.internet2.edu (64.57.28.57) 31.800 ms 41.967 ms 31.787 ms
7 so-3-0-0.0.rtr.losa.net.internet2.edu (64.57.28.44) 63.478 ms 63.704 ms 63.710 ms
8 hpr-lax-hpr--i2-newnet.cenic.net (137.164.26.132) 63.093 ms 63.026 ms 63.384 ms
9 svl-hpr--lax-hpr-10ge.cenic.net (137.164.25.13) 71.242 ms 71.542 ms 76.282 ms
10 oak-hpr--svl-hpr-10ge.cenic.net (137.164.25.9) 72.744 ms 72.243 ms 72.556 ms
11 hpr-stan-ge--oak-hpr.cenic.net (137.164.27.158) 73.763 ms 73.396 ms 73.665 ms
12 bbra-rtr.Stanford.EDU (171.64.1.134) 73.577 ms 73.682 ms 73.492 ms
13 ***
14 www5.Stanford.EDU (171.67.20.37) 77.317 ms 77.128 ms 77.648 ms
```

# Traceroute

```
$ traceroute www.msu.edu
```

```
traceroute to www.msu.edu (35.8.10.30), 64 hops max, 40 byte packets
```

```
 1 141.211.203.252 (141.211.203.252) 2.644 ms 0.973 ms 14.162 ms
 2 v-bin-seb.r-bin-seb.umnet.umich.edu (192.122.183.61) 1.847 ms 0.561 ms 0.496 ms
 3 v-bin-seb-i2-aa.merit-aa2.umnet.umich.edu (192.12.80.33) 6.490 ms 6.499 ms 6.529 ms
 4 lt-0-3-0x1.eq-chi2.mich.net (198.108.23.121) 8.096 ms 8.113 ms 8.103 ms
 5 xe-0-0-0x23.msu6.mich.net (198.108.23.213) 7.831 ms 7.962 ms 7.965 ms
 6 192.122.183.227 (192.122.183.227) 12.953 ms 12.339 ms 10.322 ms
 7 cc-tl-gel-23.net.msu.edu (35.9.101.209) 9.522 ms 9.406 ms 9.817 ms
 8 * * *
```

# Traceroute

\$ `traceroute www.pku.edu.cn`

traceroute: Warning: www.pku.edu.cn has multiple addresses; using 162.105.129.104

traceroute to www.pku.edu.cn (162.105.129.104), 64 hops max, 40 byte packets

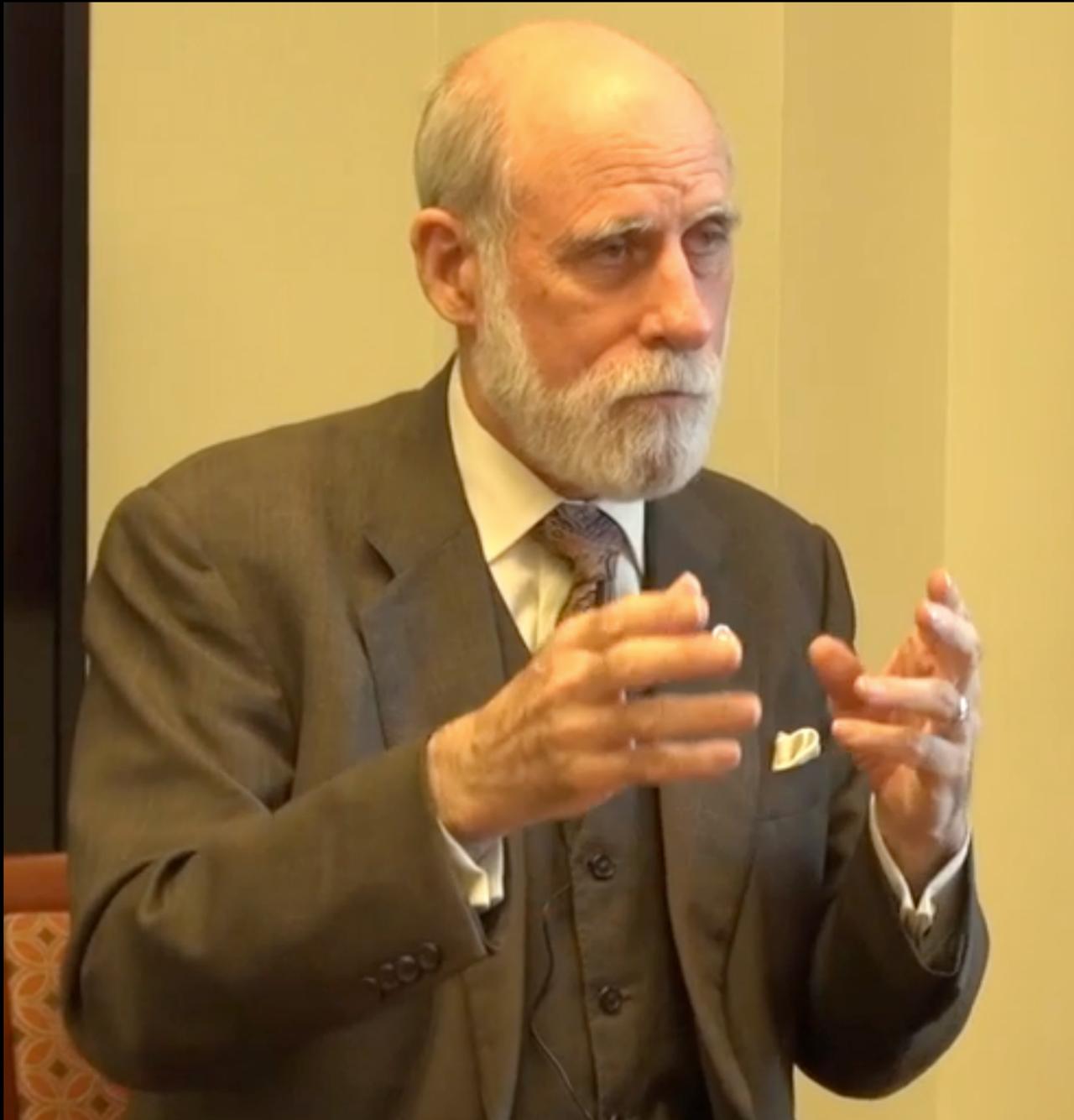
1	141.211.203.252 (141.211.203.252)	1.228 ms	0.584 ms	0.592 ms	
2	v-bin-seb.r-bin-seb.umnet.umich.edu (192.122.183.61)	0.604 ms	0.565 ms	0.466 ms	
3	v-bin-seb-i2-aa.merit-aa2.umnet.umich.edu (192.12.80.33)	7.511 ms	6.641 ms	6.588 ms	
4	192.122.183.30 (192.122.183.30)	12.078 ms	6.989 ms	7.619 ms	Michigan Tennessee
5	192.31.99.133 (192.31.99.133)	7.666 ms	8.953 ms	17.861 ms	
6	192.31.99.170 (192.31.99.170)	59.275 ms	59.273 ms	59.108 ms	
7	134.75.108.209 (134.75.108.209)	173.614 ms	173.552 ms	173.333 ms	
8	134.75.107.10 (134.75.107.10)	256.760 ms	134.75.107.18 (134.75.107.18)	256.574 ms	256.53
9	202.112.53.17 (202.112.53.17)	256.761 ms	256.801 ms	256.688 ms	Seoul
10	202.112.61.157 (202.112.61.157)	257.416 ms	257.960 ms	257.747 ms	
11	202.112.53.194 (202.112.53.194)	256.827 ms	257.068 ms	256.962 ms	
12	202.112.41.202 (202.112.41.202)	256.800 ms	257.053 ms	256.933 ms	Beijing

# The perfect is the enemy of the good

*Le mieux est l'ennemi du bien. --Voltaire*

- **IP Does:** Best effort to get data across bunch of hops from one network to another network
- **IP Does Not:** Guarantee delivery - if things go bad - the data can vanish
- Best effort to keep track of the good and bad paths for traffic - tries to pick better paths when possible
- This makes it fast and scalable to very large networks - and ultimately “reliable” because it does not try to do too much

# Vint Cerf: A Brief History of Packets



- Instrumental in the design and development of the ARPANET
- Vint was a graduate student as the notions of packet-switching were emerging across academia

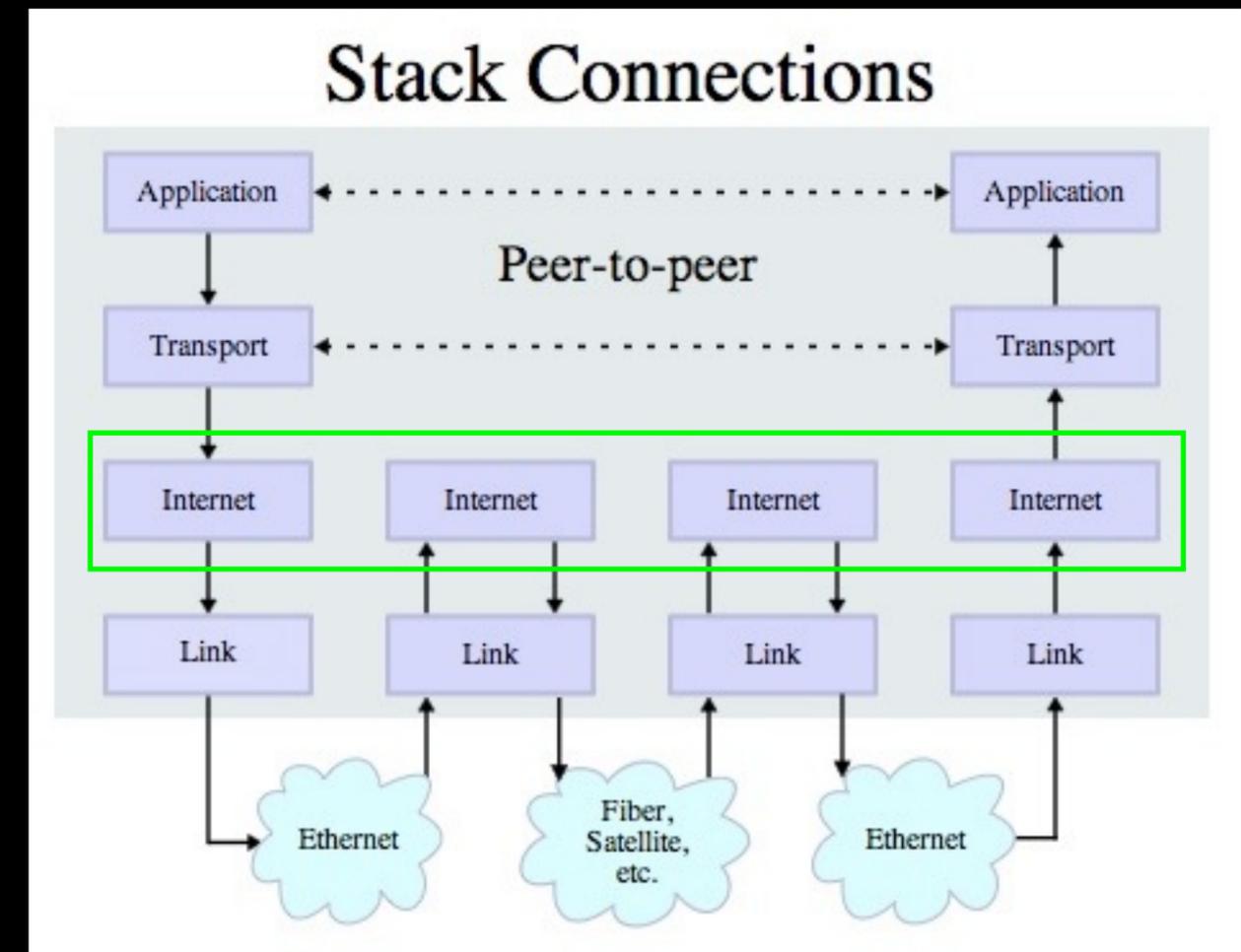
# Domain Name System

**The Domain Name System  
convert user-friendly  
names, like**

**www.umich.edu**

**to network-friendly IP  
addresses, like**

**141.211.32.166**



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Domain Name System

- Numeric addresses like 141.211.63.45 are great for Internet routers but lousy for people
- Each campus ends up with a lot of networks (141.211.\*, 65.43.21.\*)
- Sometimes (rarely) the IP address numbers get reorganized
- When servers physically move they need new IP addresses

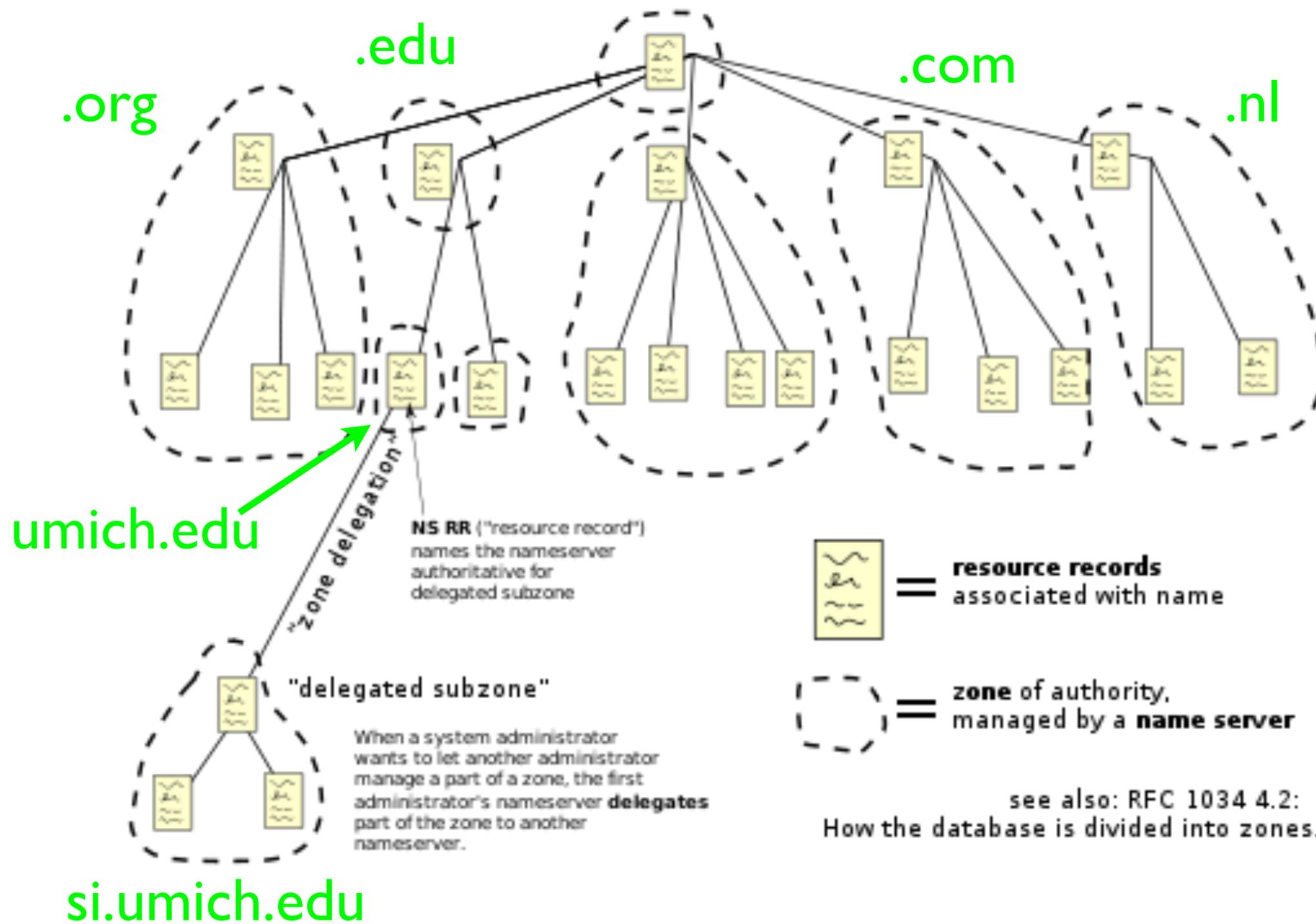
# DNS: Internet Address Book

- The Domain Name System is a big fast distributed database of Internet names to Internet “phone numbers”
- IP Addresses reflect technical “geography”
  - **141.211.63.44** - read left to right like a phone number  

- Domain names reflect organizational structure
  - **www.si.umich.edu** - read right to left like postal address  

  - **2455 North Quad, Ann Arbor, MI, USA, Earth**  


# Domain Name Space



# Internetwork Layer (IP)

[http://en.wikipedia.org/wiki/Internet\\_Protocol](http://en.wikipedia.org/wiki/Internet_Protocol)

<http://en.wikipedia.org/wiki/Traceroute>

<http://en.wikipedia.org/wiki/Ping>

Application Layer  
Web, E-Mail, File Transfer

Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Scalable, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Transport Layer

[http://en.wikipedia.org/wiki/  
Transmission\\_Control\\_Protocol](http://en.wikipedia.org/wiki/Transmission_Control_Protocol)

Application Layer  
Web, E-Mail, File Transfer

Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Scalable, Unreliable

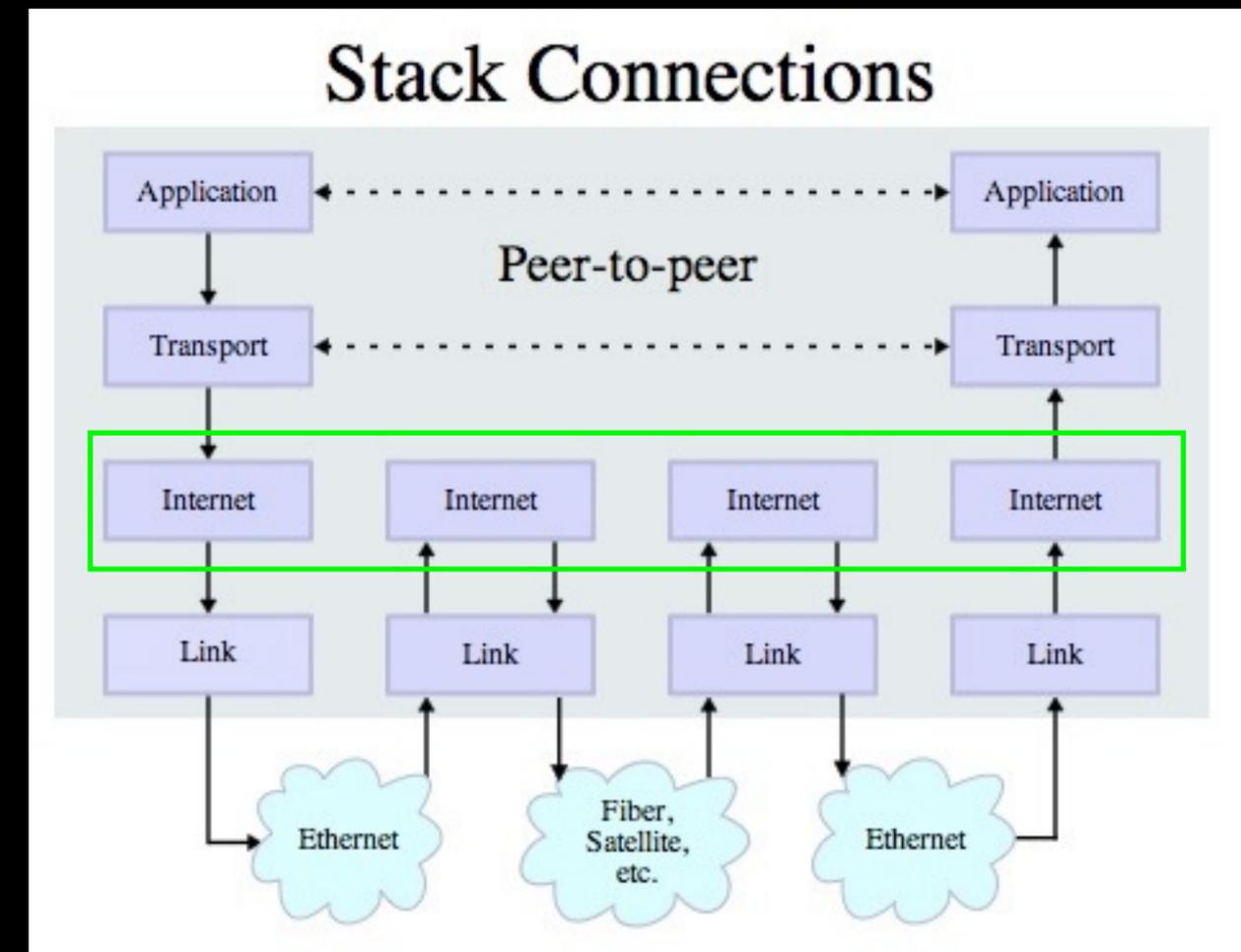
Link Layer (Ethernet, WiFi)  
Physical Connections

# Review: The Magic of IP

- What it does - Tries to get one packet across a 5-20 of hops from one network to another network
- Keeps track of the good and bad paths for traffic - tries to pick better paths when possible
- But no guarantee of delivery - if things go bad - the data vanishes
- This makes it fast and scalable - and ultimately “reliable” because it does not try to do too "everything"

# Internet Protocol

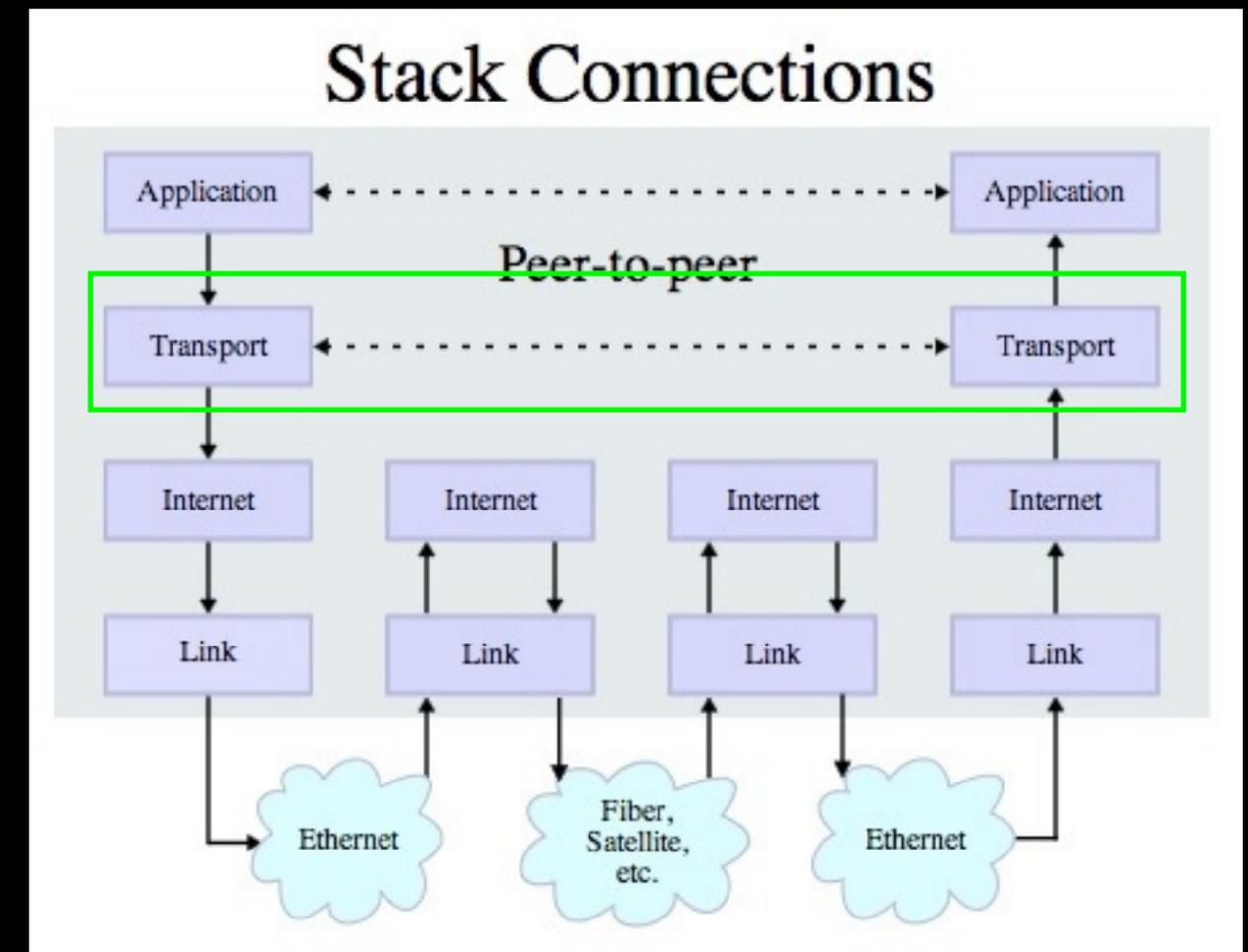
- So many links / hops
- So many routes
- Things can change dynamically and IP has to react (links up/down)
- IP can drop packets



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Transmission Protocol (TCP)

- Built on top of IP
- Assumes IP might lose some data
- In case data gets lost - **we keep a copy of the data** and we send until we get an acknowledgement
- If it takes “too long” - just send it again



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

Sender

100

200

300

400

500

Break Messages  
into Pieces

Receiver

Sender

100

200

300

400

500

100

200

300

Receiver

Break Messages  
into Pieces

Sender

100



200

300

400

500

Got 100  
Where is 200

Receiver

100

300

Break Messages  
into Pieces

Sender

100



200



300

400

500

Got 200

Receiver

100

Break Messages  
into Pieces

Sender

300



400



500

Got 400

Receiver

100

200

300

400

500

Break Messages  
into Pieces

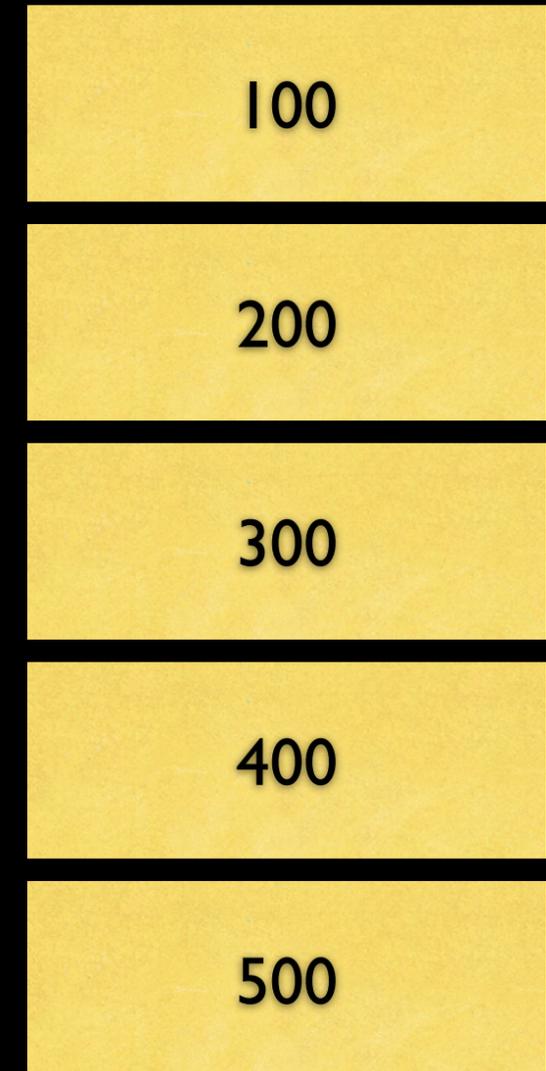
Sender



Break Messages  
into Pieces

Got 500

Receiver



Sender

Receiver

100

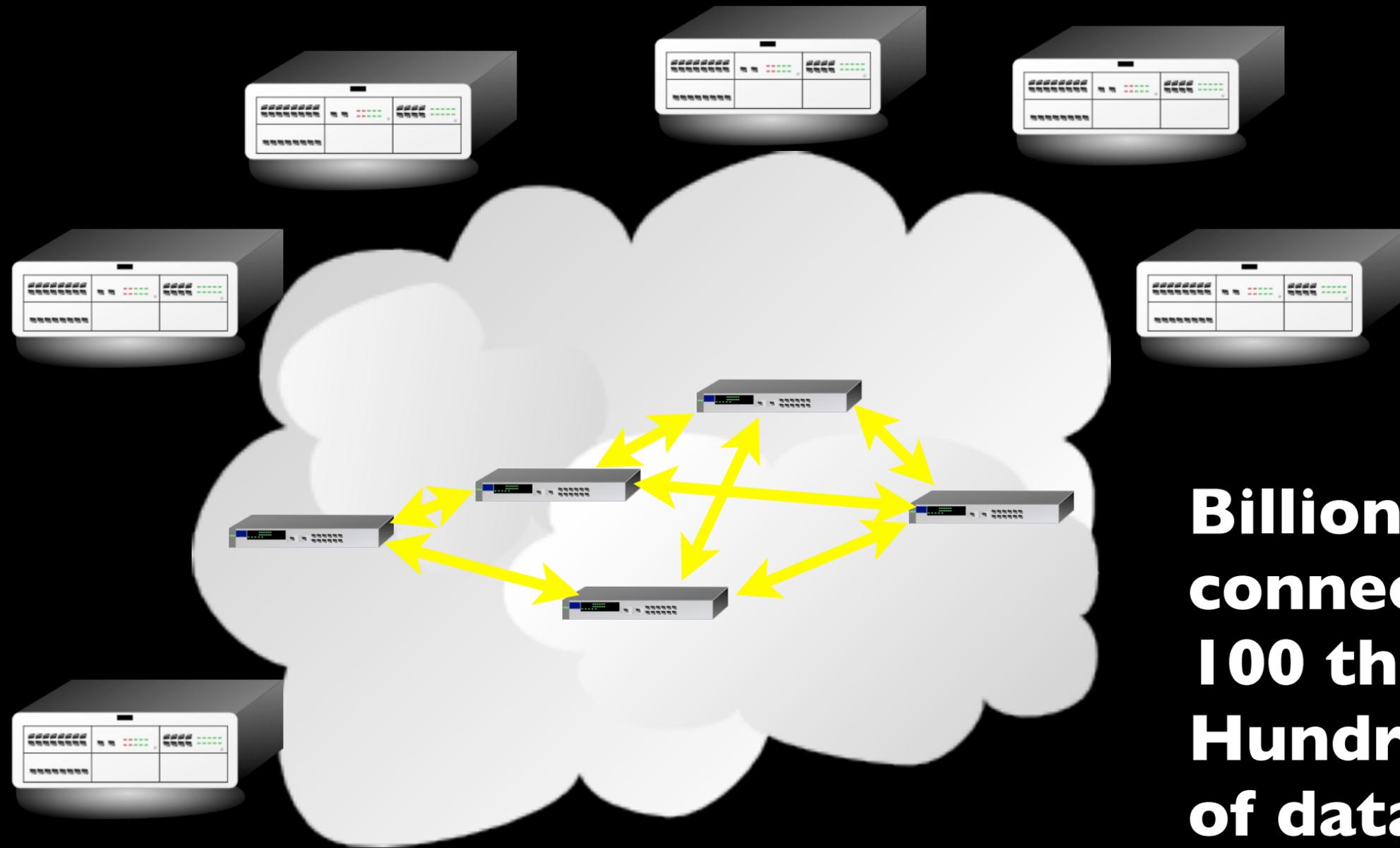
200

300

400

500

Break Messages  
into Pieces



**Billions of computers connected to the internet; 100 thousands of routers. Hundreds of billions bytes of data enroute at any moment.**

**Storage of enroute data done at the edges only!**

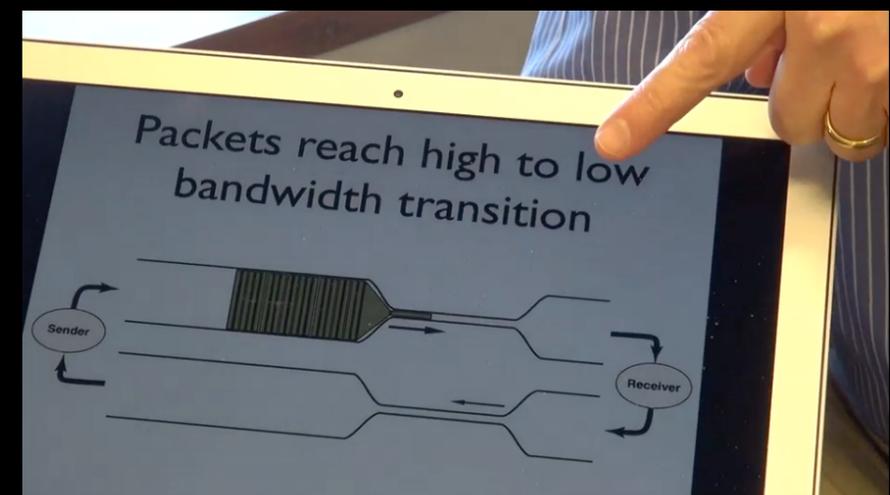
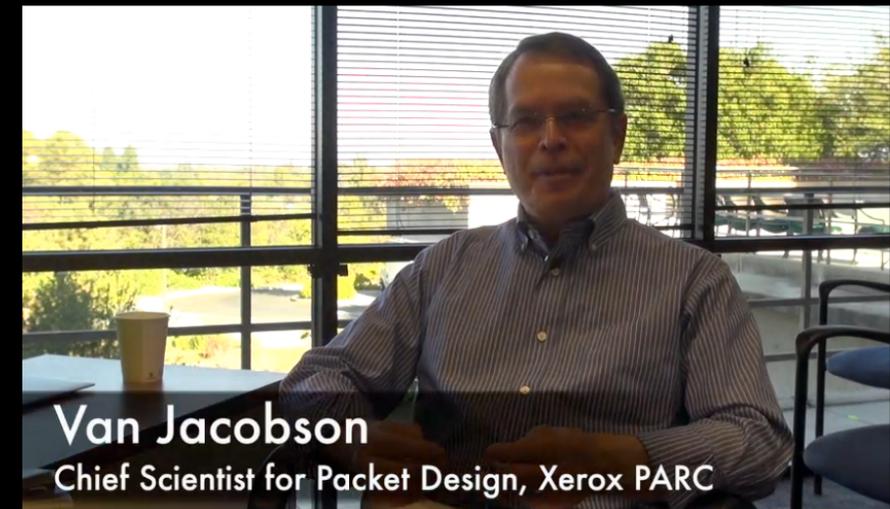
# One (of many) Scary Problem(s)

- In 1987 as local campuses with 10 MBit networks were connected together using 56Kbit leased lines, things kind of fell apart
- At some point, when there was a little too much traffic, it all fell apart...

<http://www.youtube.com/watch?v=IVgIMeRYmWI>

[http://en.wikipedia.org/wiki/Van\\_Jacobson](http://en.wikipedia.org/wiki/Van_Jacobson)

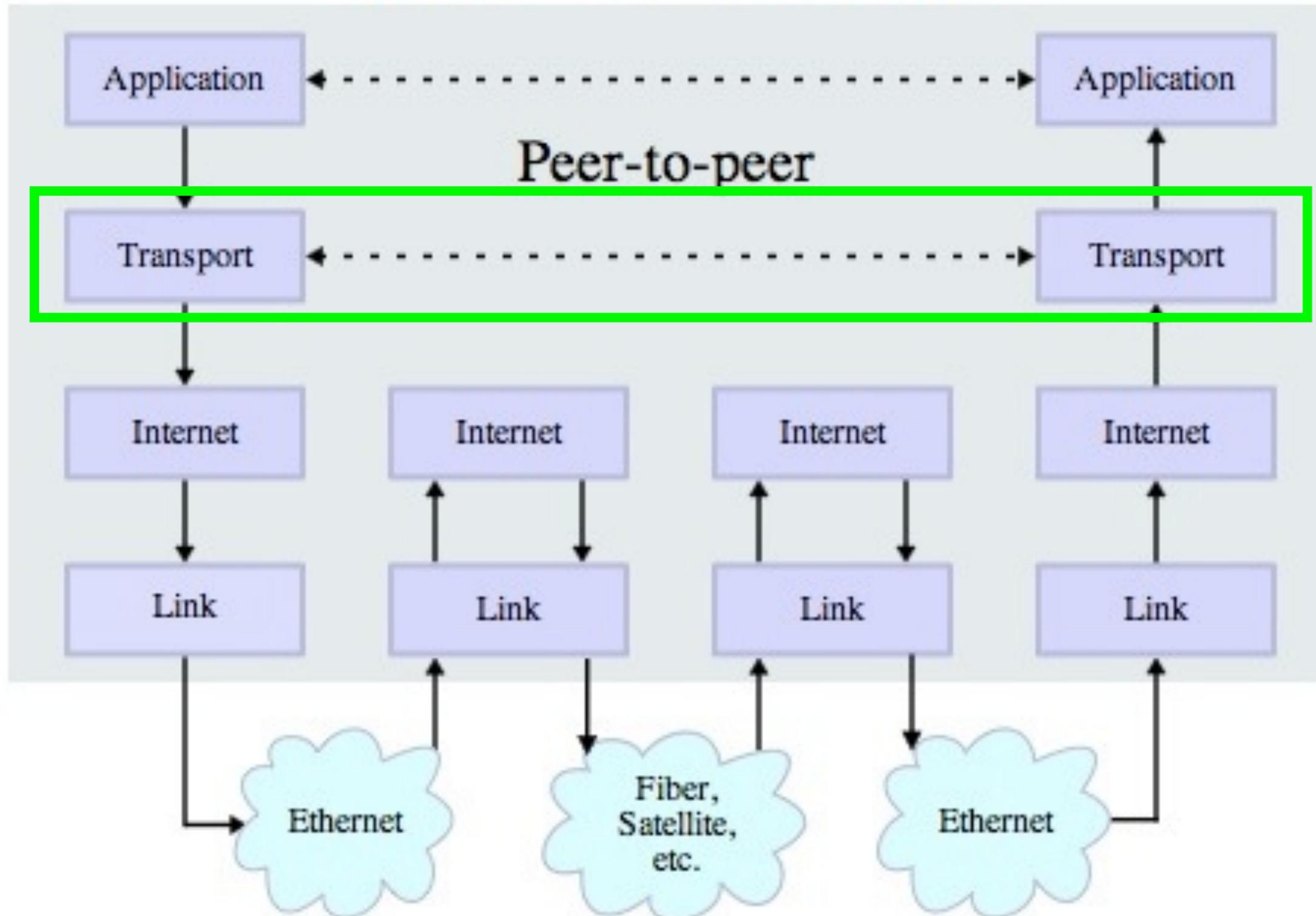
[http://en.wikipedia.org/wiki/TCP\\_congestion\\_avoidance\\_algorithm](http://en.wikipedia.org/wiki/TCP_congestion_avoidance_algorithm)



# Transmission Protocol (TCP)

- The responsibility of the transport layer is to present a reliable end-to-end pipe to the application
- Data either arrives in the proper order or the connection is closed
- TCP keeps buffers in the sending and destination system to keep data which has arrived out of order or to retransmit if necessary
- TCP provides individual connections between applications

# Stack Connections



# Application Layer

Application Layer  
Web, E-Mail, File Transfer

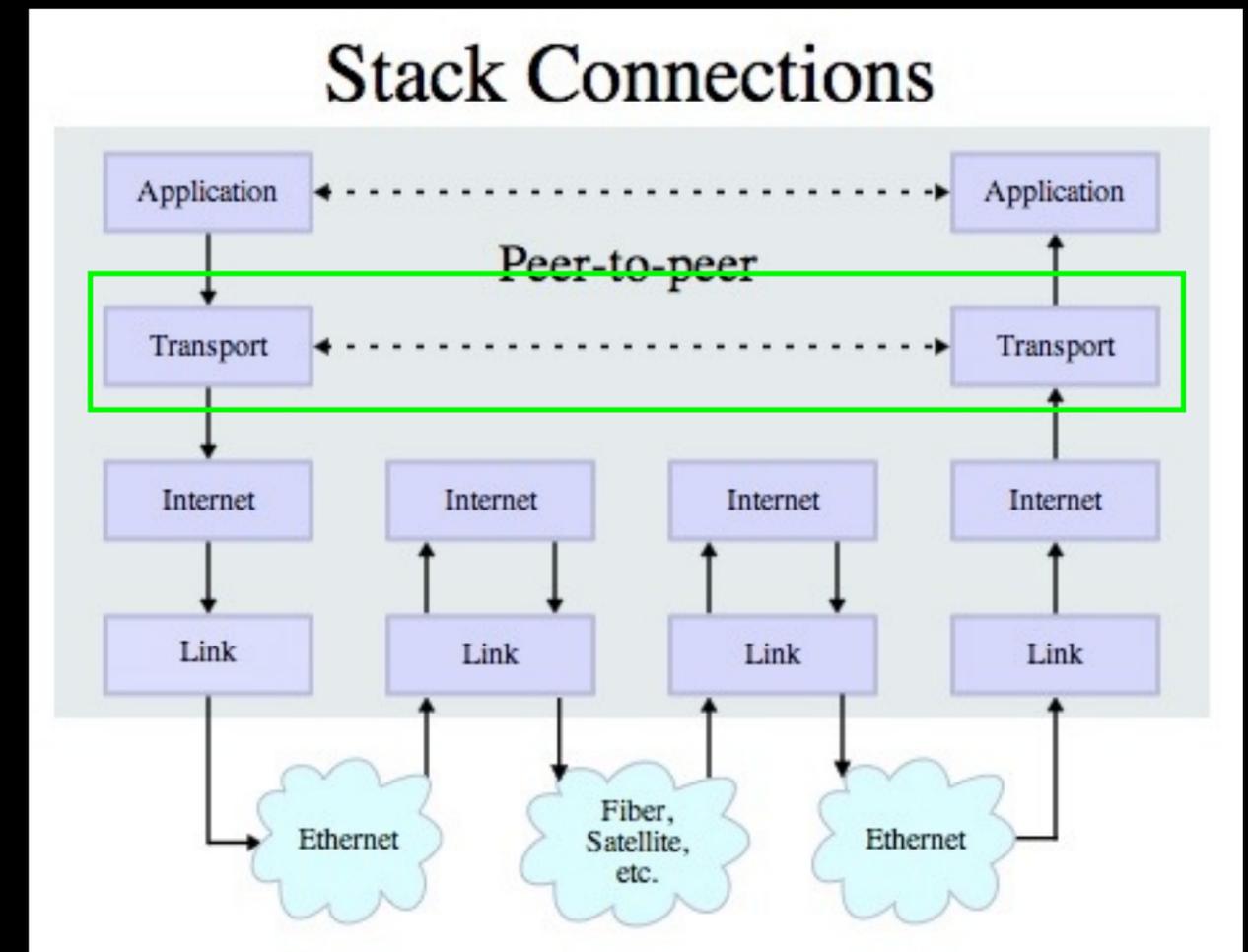
Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# Quick Review

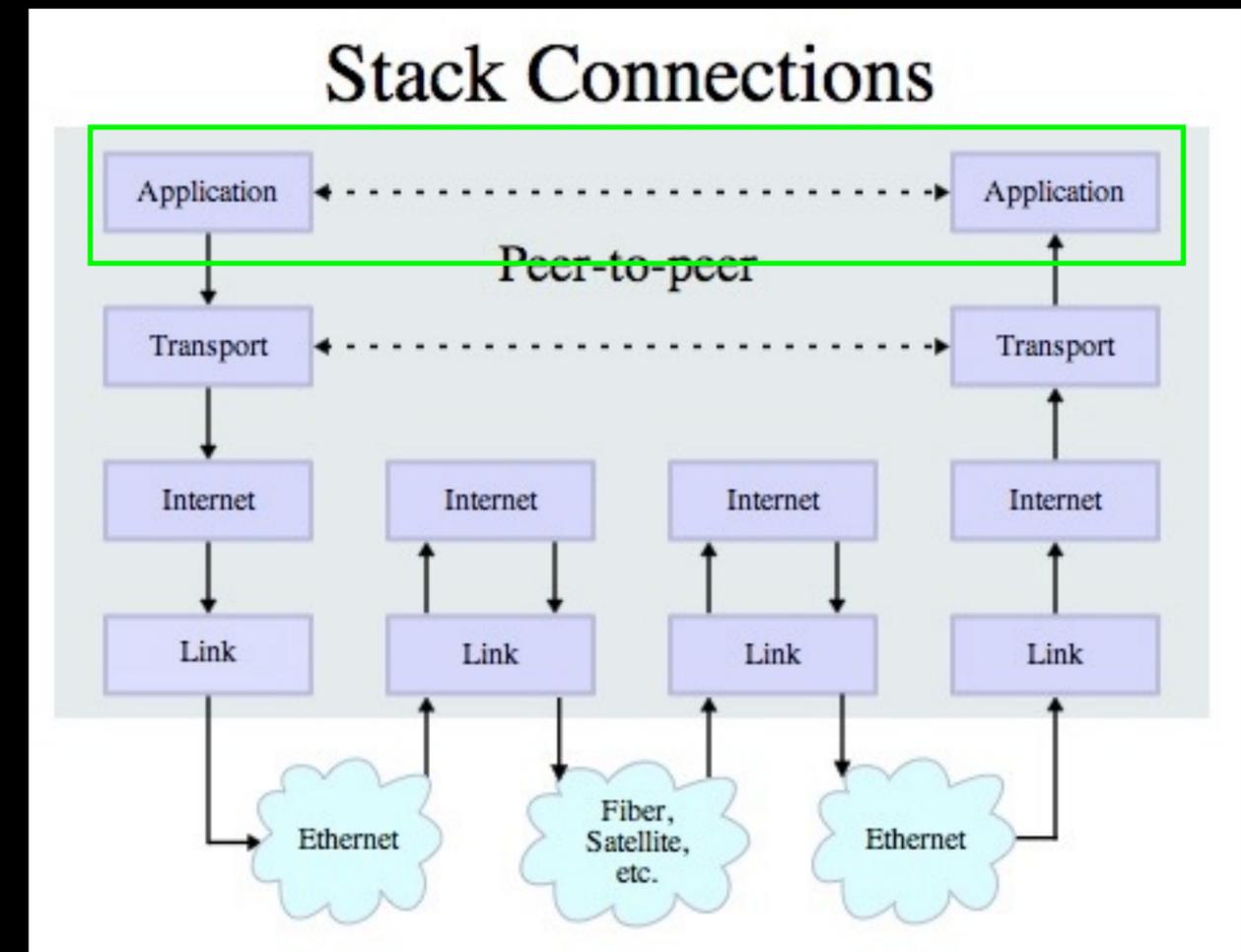
- Link layer: gets the data onto the link, and manages collisions on a single hop
- Internet layer: moves the data over one hop, trying to get it “closer” to its destination
- Transport layer: Assumes that the internet layer may lose data, so request retransmission when needed— provides a nice reliable pipe from source to destination



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Application Protocol

- Since TCP gives us a reliable pipe, what do we want to do with the pipe? What problem do we want to solve?
  - Mail
  - World Wide Web
  - Stream kitty videos



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# Two Questions for the Application Layer

- Which application gets the data?
  - Ports
- What are the rules for talking with that application?
  - Protocols

[http://en.wikipedia.org/wiki/TCP\\_and\\_UDP\\_port](http://en.wikipedia.org/wiki/TCP_and_UDP_port)

[http://en.wikipedia.org/wiki/List\\_of\\_TCP\\_and\\_UDP\\_port\\_numbers](http://en.wikipedia.org/wiki/List_of_TCP_and_UDP_port_numbers)

# Ports

- Like extensions in a phone number
- The IP address network number (the area code) gets to the LAN
- The IP address host number (the telephone number) gets you to the destination machine
- The port (the extension) gets you to a specific application

(734) 764 1855, ext. 27

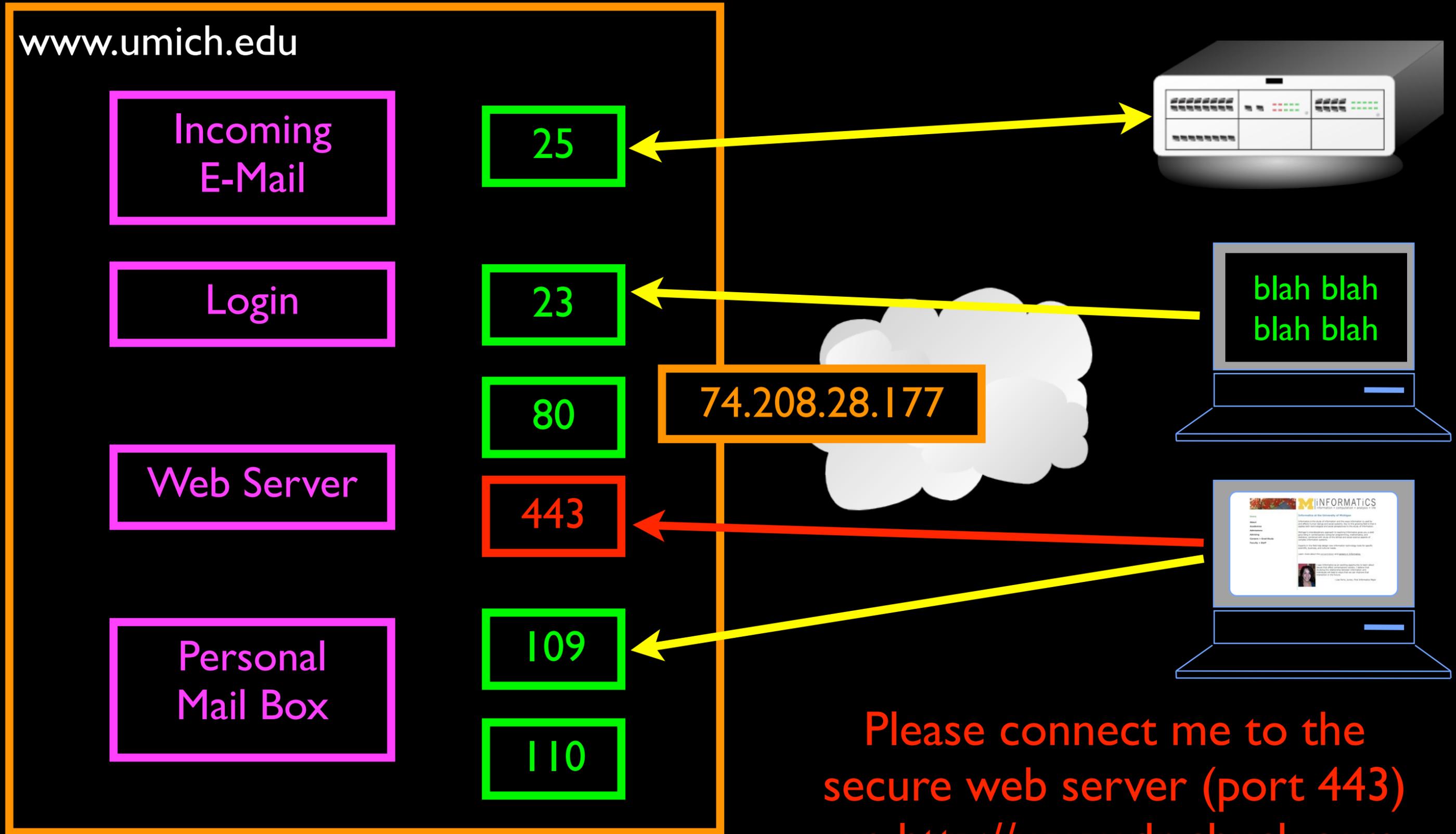
141.211.144.188

Port 25

# TCP, Ports, and Connections

[http://en.wikipedia.org/wiki/TCP\\_and\\_UDP\\_port](http://en.wikipedia.org/wiki/TCP_and_UDP_port)

[http://en.wikipedia.org/wiki/List\\_of\\_TCP\\_and\\_UDP\\_port\\_numbers](http://en.wikipedia.org/wiki/List_of_TCP_and_UDP_port_numbers)



Please connect me to the secure web server (port 443) on <http://www.dr-chuck.com>

# Common TCP Ports

- Telnet (23) - Login
- SSH (22) - Secure Login
- HTTP (80)
- HTTPS (443) - Secure
- SMTP (25) (Mail)
- IMAP (143/220/993) - Mail Retrieval
- POP (109/110) - Mail Retrieval
- DNS (53) - Domain Name
- FTP (21) - File Transfer

[http://en.wikipedia.org/wiki/List\\_of\\_TCP\\_and\\_UDP\\_port\\_numbers](http://en.wikipedia.org/wiki/List_of_TCP_and_UDP_port_numbers)

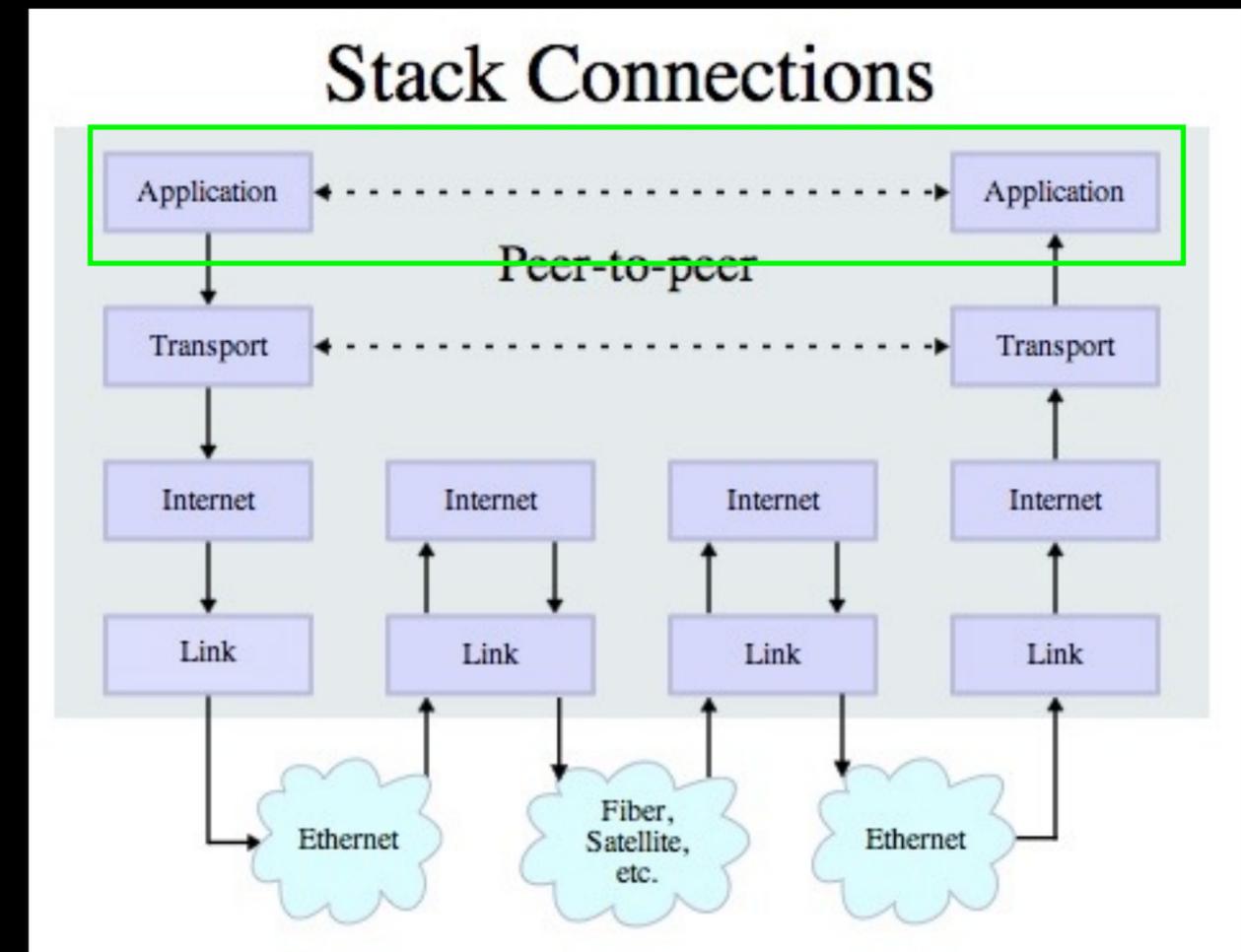
# Application Protocols

<http://en.wikipedia.org/wiki/Http>

<http://en.wikipedia.org/wiki/Pop3>

# Application Protocol

- Since TCP gives us a reliable pipe, what to we want to do with the pipe? What problem do we want to solve?
  - Mail
  - World Wide Web
  - Stream kitty videos



Source: [http://en.wikipedia.org/wiki/Internet\\_Protocol\\_Suite](http://en.wikipedia.org/wiki/Internet_Protocol_Suite)

# HTTP - Hypertext Transport Protocol

- The dominant Application Layer Protocol on the Internet
- Invented for the Web - to Retrieve HTML, Images, Documents etc
- Extended to be data in addition to documents - RSS, Web Services, etc..
- Basic Concept - Make a Connection - Request a document - Retrieve the Document - Close the Connection

<http://en.wikipedia.org/wiki/Http>

# HTTP Request / Response Cycle

Web Server

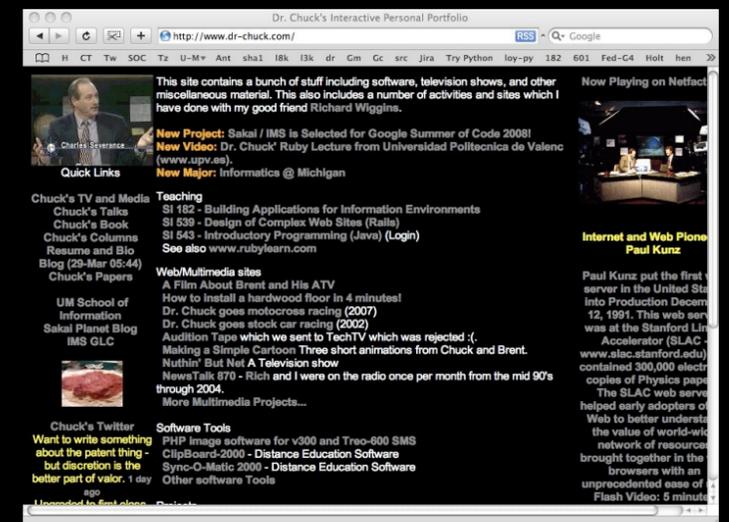
HTTP Request

HTTP Response

Browser

Hello there my name is Chuck  
Go ahead and click on [here](#).

Internet Explorer,  
FireFox, Safari, etc.



[http://www.oreilly.com/openbook/cgi/ch04\\_02.html](http://www.oreilly.com/openbook/cgi/ch04_02.html)

Source: <http://www.dr-chuck.com/>

# HTTP Request / Response Cycle

Web Server

```
<head> .. </head>
<body>
<h1>Welcome to my
application</h1>
...
</body>
```

GET /page2.html

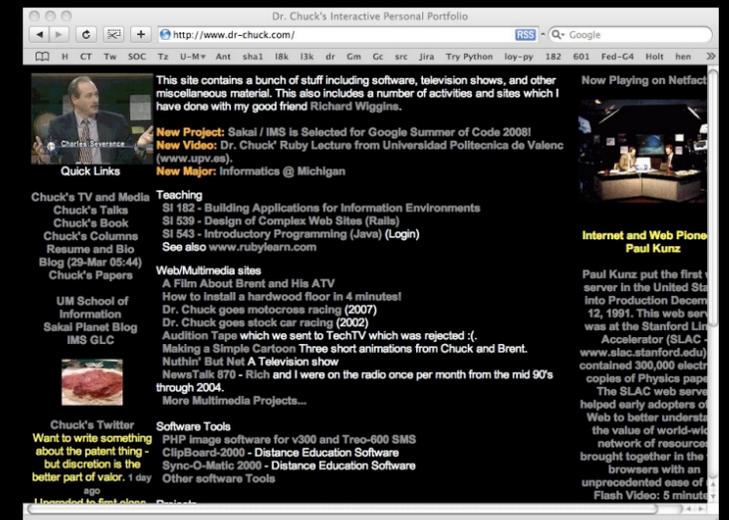
HTTP Request

HTTP Response

Browser

Hello there my name is Chuck  
Go ahead and click on [here](#).

Internet Explorer,  
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[http://www.oreilly.com/openbook/cgi/ch04\\_02.html](http://www.oreilly.com/openbook/cgi/ch04_02.html)

Source: <http://www.dr-chuck.com/>

# Internet Standards

- The standards for all of the Internet protocols (inner workings) are developed by an organization
- Internet Engineering Task Force (IETF)
- [www.ietf.org](http://www.ietf.org)
- Standards are called “RFCs” - “Request for Comments”

```
Network Working Group                                T. Berners-Lee
Request for Comments: 1945                            MIT/LCS
Category: Informational                               R. Fielding
                                                       UC Irvine
                                                       H. Frystyk
                                                       MIT/LCS
                                                       May 1996

                Hypertext Transfer Protocol -- HTTP/1.0

Status of This Memo

This memo provides information for the Internet community. This memo
does not specify an Internet standard of any kind. Distribution of
this memo is unlimited.

IESG Note:

The IESG has concerns about this protocol, and expects this document
to be replaced relatively soon by a standards track document.

Abstract

The Hypertext Transfer Protocol (HTTP) is an application-level
protocol with the lightness and speed necessary for distributed,
collaborative, hypermedia information systems. It is a generic,
stateless, object-oriented protocol which can be used for many tasks,
such as name servers and distributed object management systems,
through extension of its request methods (commands). A feature of
HTTP is the typing of data representation, allowing systems to be
built independently of the data being transferred.
```

Source: <http://www.ietf.org/rfc/rfc1945.txt>

### 5.1.2 Request-URI

The Request-URI is a Uniform Resource Identifier (Section 3.2) and identifies the resource upon which to apply the request.

$$\text{Request-URI} = \text{absoluteURI} \mid \text{abs\_path}$$

The two options for Request-URI are dependent on the nature of the request.

The absoluteURI form is only allowed when the request is being made to a proxy. The proxy is requested to forward the request and return the response. If the request is GET or HEAD and a prior response is cached, the proxy may use the cached message if it passes any restrictions in the Expires header field. Note that the proxy may forward the request on to another proxy or directly to the server specified by the absoluteURI. In order to avoid request loops, a proxy must be able to recognize all of its server names, including any aliases, local variations, and the numeric IP address. An example Request-Line would be:

```
GET http://www.w3.org/pub/WWW/TheProject.html HTTP/1.0
```

The most common form of Request-URI is that used to identify a resource on an origin server or gateway. In this case, only the absolute path of the URI is transmitted (see Section 3.2.1, `abs_path`). For example, a client wishing to retrieve the resource above directly from the origin server would create a TCP connection to port 80 of the host "www.w3.org" and send the line:

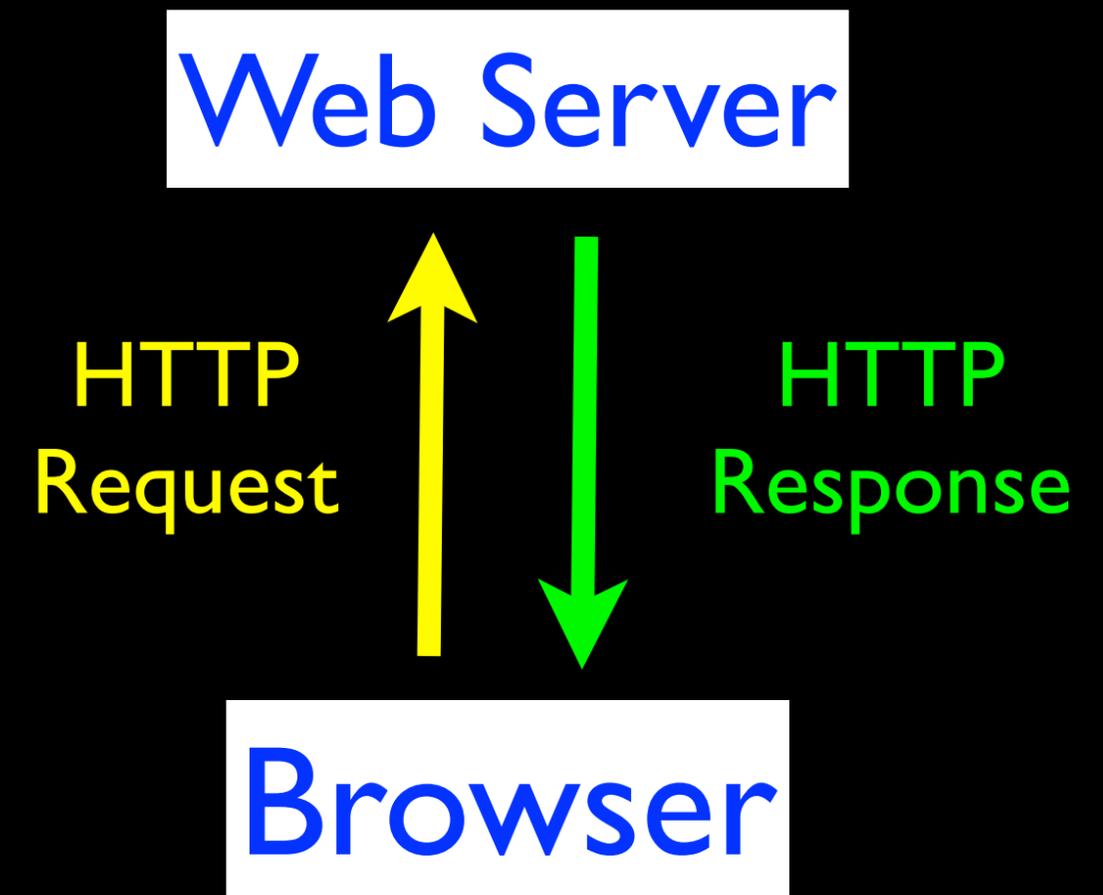
```
GET /pub/WWW/TheProject.html HTTP/1.0
```

followed by the remainder of the Full-Request. Note that the absolute path cannot be empty; if none is present in the original URI, it must be given as "/" (the server root).

The Request-URI is transmitted as an encoded string, where some characters may be escaped using the "% HEX HEX" encoding defined by RFC 1738 [4]. The origin server must decode the Request-URI in order to properly interpret the request.

# “Hacking” HTTP

```
Last login: Wed Oct 10 04:20:19 on tty2
si-csev-mbp:~ csev$ telnet www.dr-chuck.com 80
Trying 74.208.28.177...
Connected to www.dr-chuck.com.
Escape character is '^]'.
GET http://www.dr-chuck.com/page1.htm
<h1>The First Page</h1>
<p>
If you like, you can switch to the
<a href="http://www.dr-chuck.com/page2.htm">
Second Page</a>.
</p>
```



Port 80 is the non-encrypted HTTP port

# Accurate Hacking in the Movies

- Matrix Reloaded
- Bourne Ultimatum
- Die Hard 4
- ...

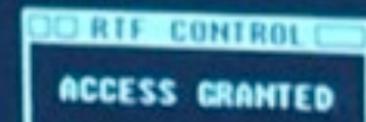


```
80/tcp    open     http
81/tcp    open     hosts2-ns
10[0]     [noble]
11 # nmap -u -sS -O 10.2.2.2
11
13 Starting nmap V. 2.54BETA25
13 Insufficient responses for TCP sequencing (3), OS detection i
13 accurate
14 Interesting ports on 10.2.2.2:
44 (The 1539 ports scanned but not shown below are in state: cl
51 Port      State      Service
51 22/tcp    open      ssh
58
68 No exact OS matches for host
68
24 Nmap run completed -- 1 IP address (1 host up) scanned
50 # sshnuke 10.2.2.2 -rootpw="210N0101"
Connecting to 10.2.2.2:ssh ... successful.
Re Attempting to exploit SSHv1 CRC32 ... successful.
IP Resetting root password to "210N0101".
System open: Access Level <9>
50 # ssh 10.2.2.2 -l root
root@10.2.2.2's password: █
```

<http://nmap.org/movies.html> (scroll down for video)  
Or search YouTube for "Trinity hacking scene"



```
80/tcp    open      http
81/tcp    open      hosts2.nc
10:~$
11:~$ nmap -v -sS -O 10.2.2.2
11:~$
13: Starting nmap U. 2.54BETA25
13: Insufficient responses for TCP sequencing (3), OS detection
13: accurate
14: Interesting ports on 10.2.2.2:
44: (The 1539 ports scanned but not shown below are in state: cl
51: Port      State      Service
51: 22/tcp    open      ssh
58:
68: No exact OS matches for host
68:
24: Nmap run completed -- 1 IP address (1 host up) scanned
50:~$ sshnuke 10.2.2.2 -rootpw="210ND101"
Connecting to 10.2.2.2:ssh ... successful.
Re Attempting to exploit SSHv1 CRC32 ... successful.
IP Resetting root password to "210ND101".
System open: Access Level <9>
50:~$ ssh 10.2.2.2 -l root
root@10.2.2.2's password: █
```

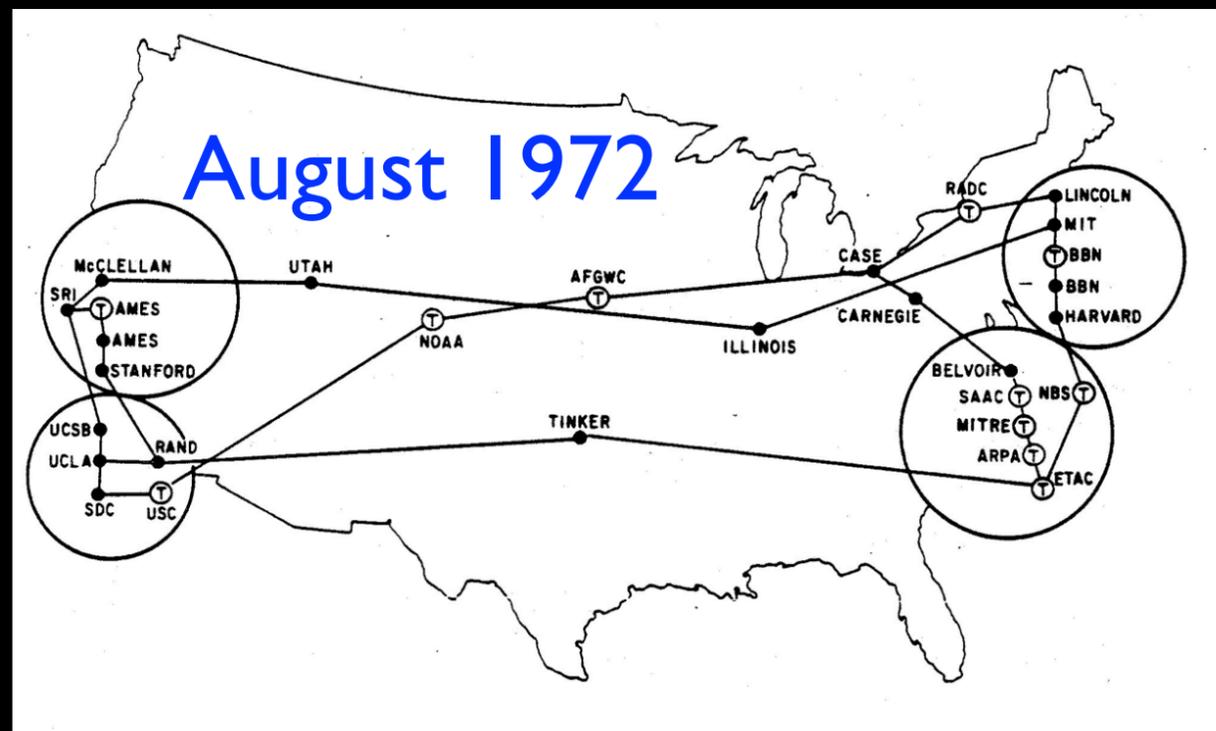


<http://nmap.org/movies.html> (scroll down for video)  
Or search YouTube for "Trinity hacking scene"

# Application Layer Summary

- We start with a “pipe” abstraction - we can send and receive data on the same “socket”
- We can optionally add a security layer to TCP using SSL - Secure Socket Layer (aka TLS - Transport Layer Security)
- We use well known “port numbers” so that applications can find a particular application *\*within\** a server such as a mail server, web service, etc

# The Architecture of the Internet



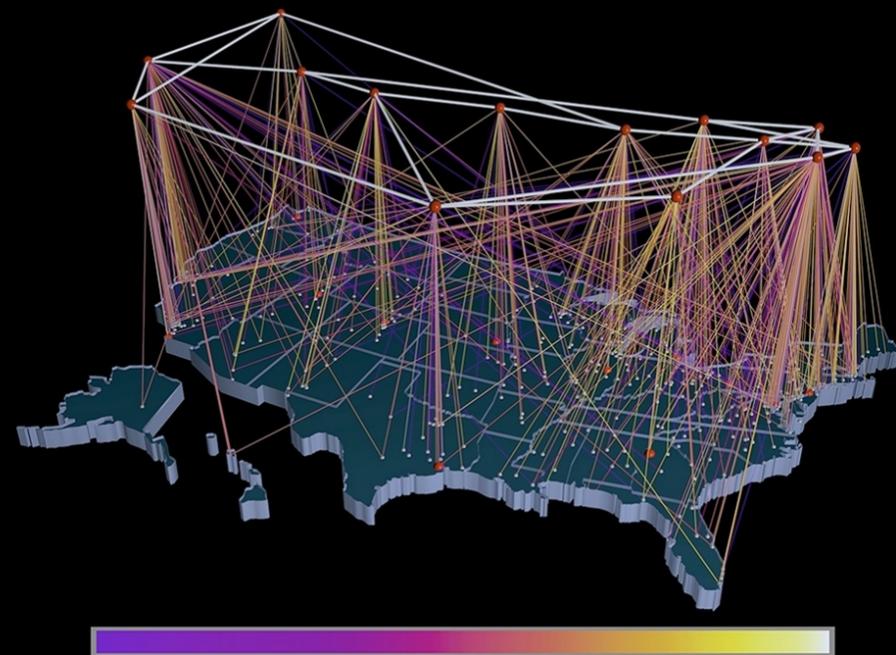
Application Layer  
Web, E-Mail, File Transfer

Transport Layer (TCP)  
Reliable Connections

Internetwork Layer (IP)  
Simple, Unreliable

Link Layer (Ethernet, WiFi)  
Physical Connections

# The Architecture of the Internet



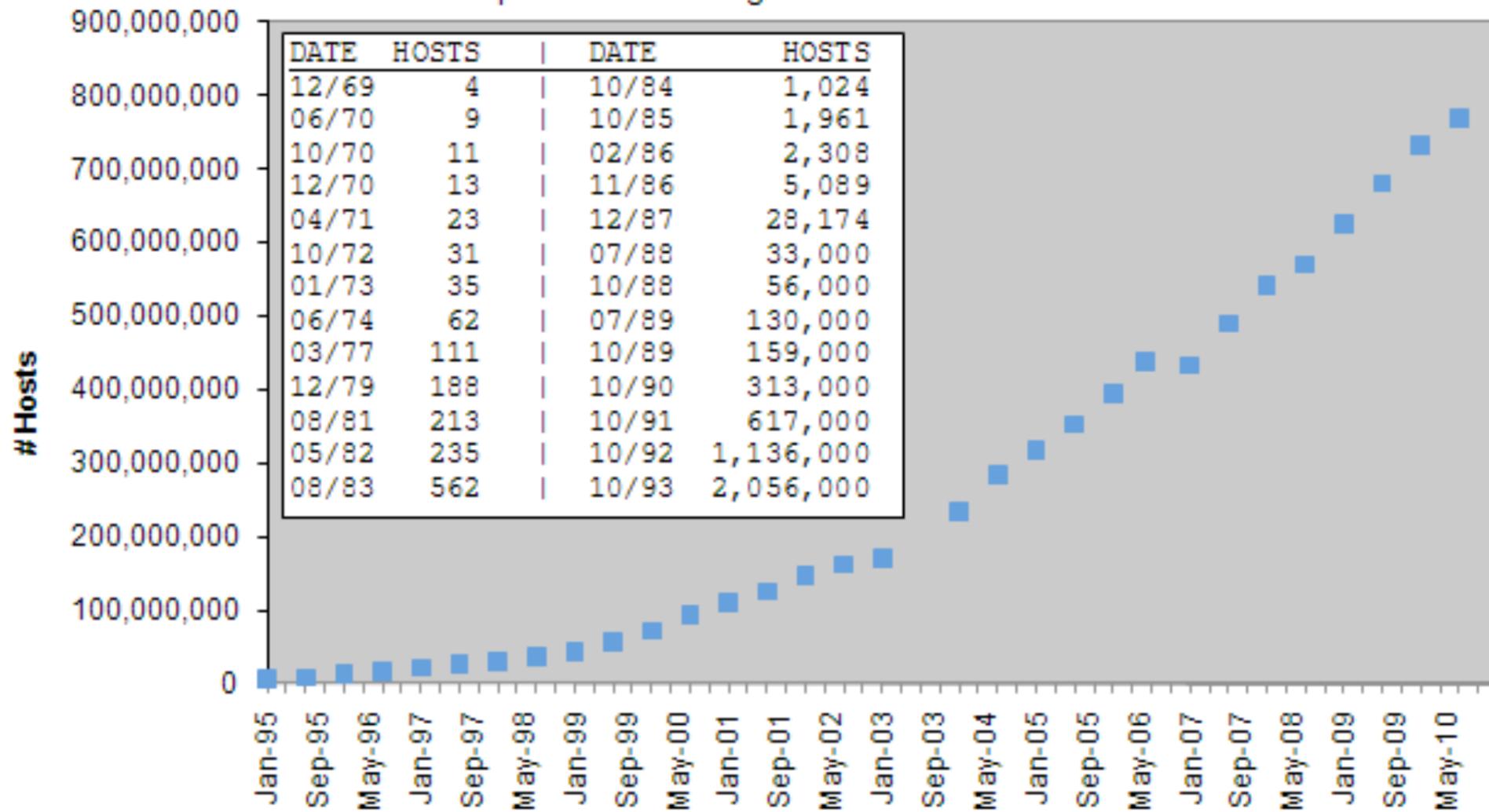
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Physical Connections

Hobbes' Internet Timeline Copyright ©2010 Robert H Zakon  
<http://www.zakon.org/robert/internet/timeline/>

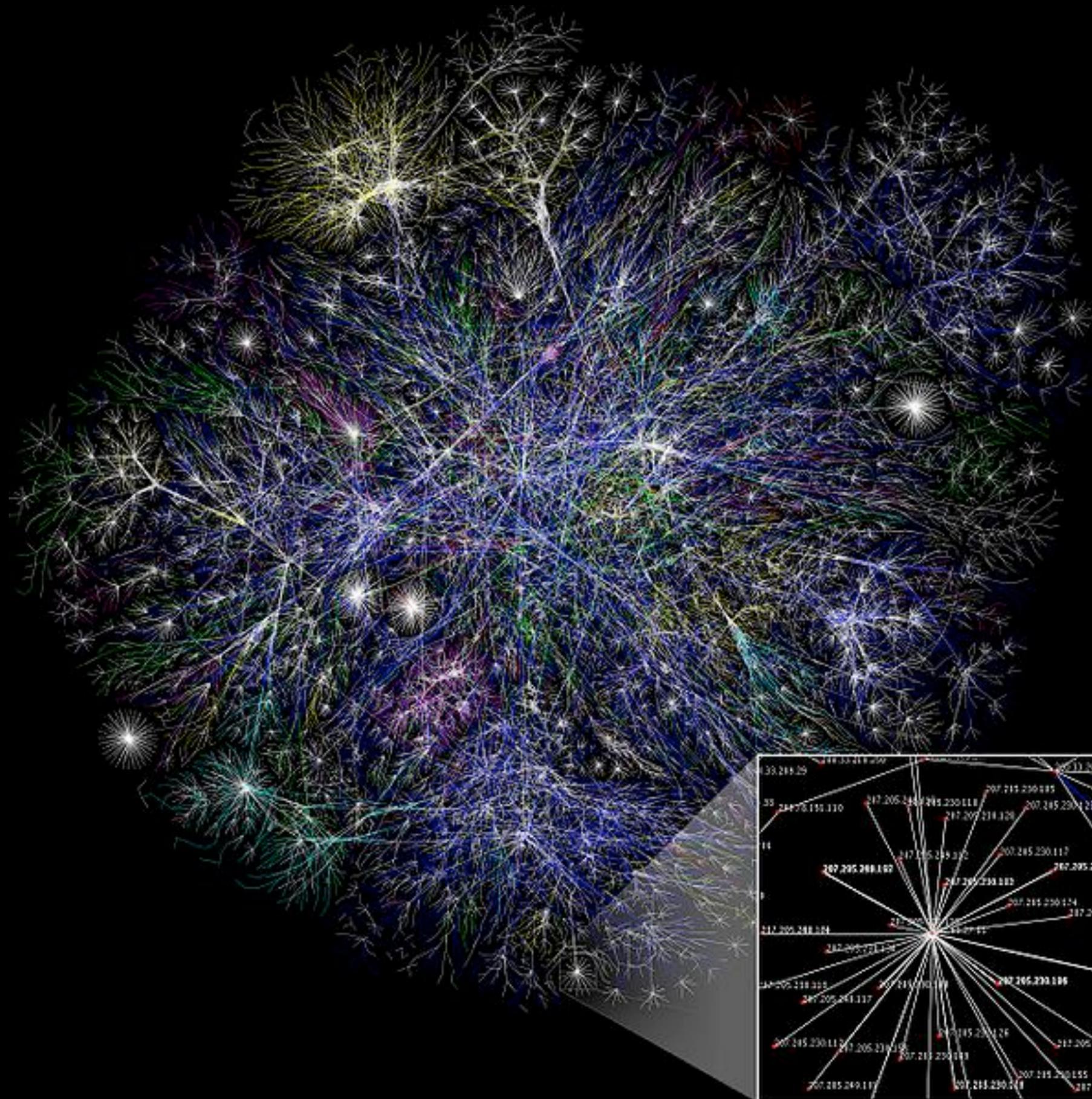


**Application Layer**  
Web, E-Mail, File Transfer

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Application Layer  
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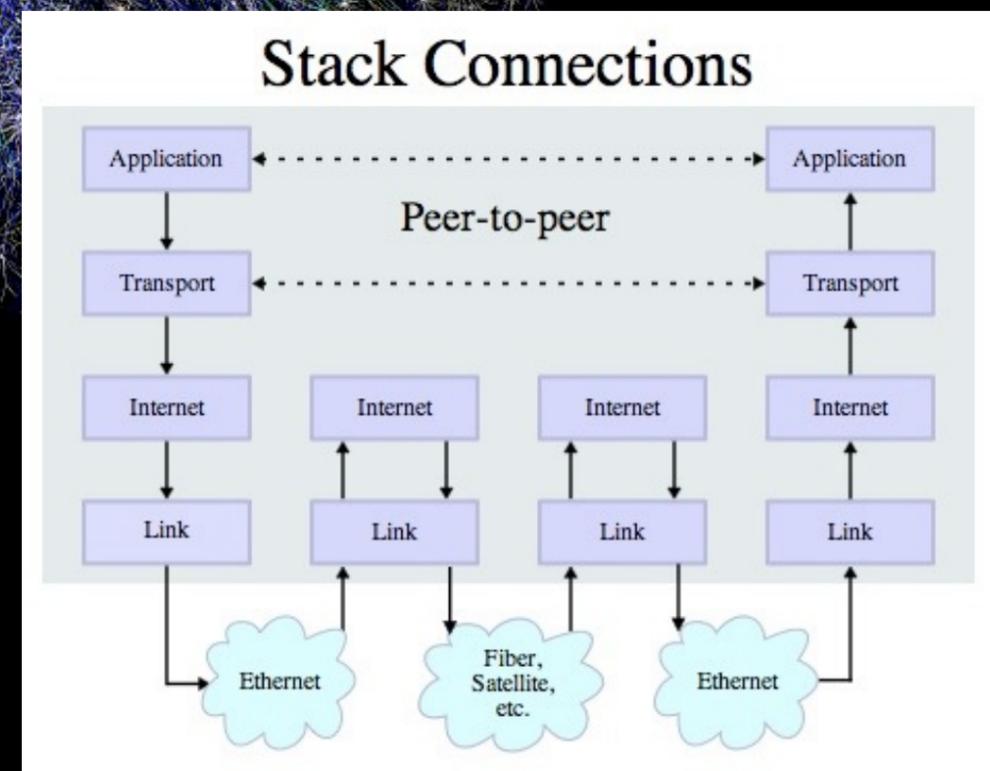
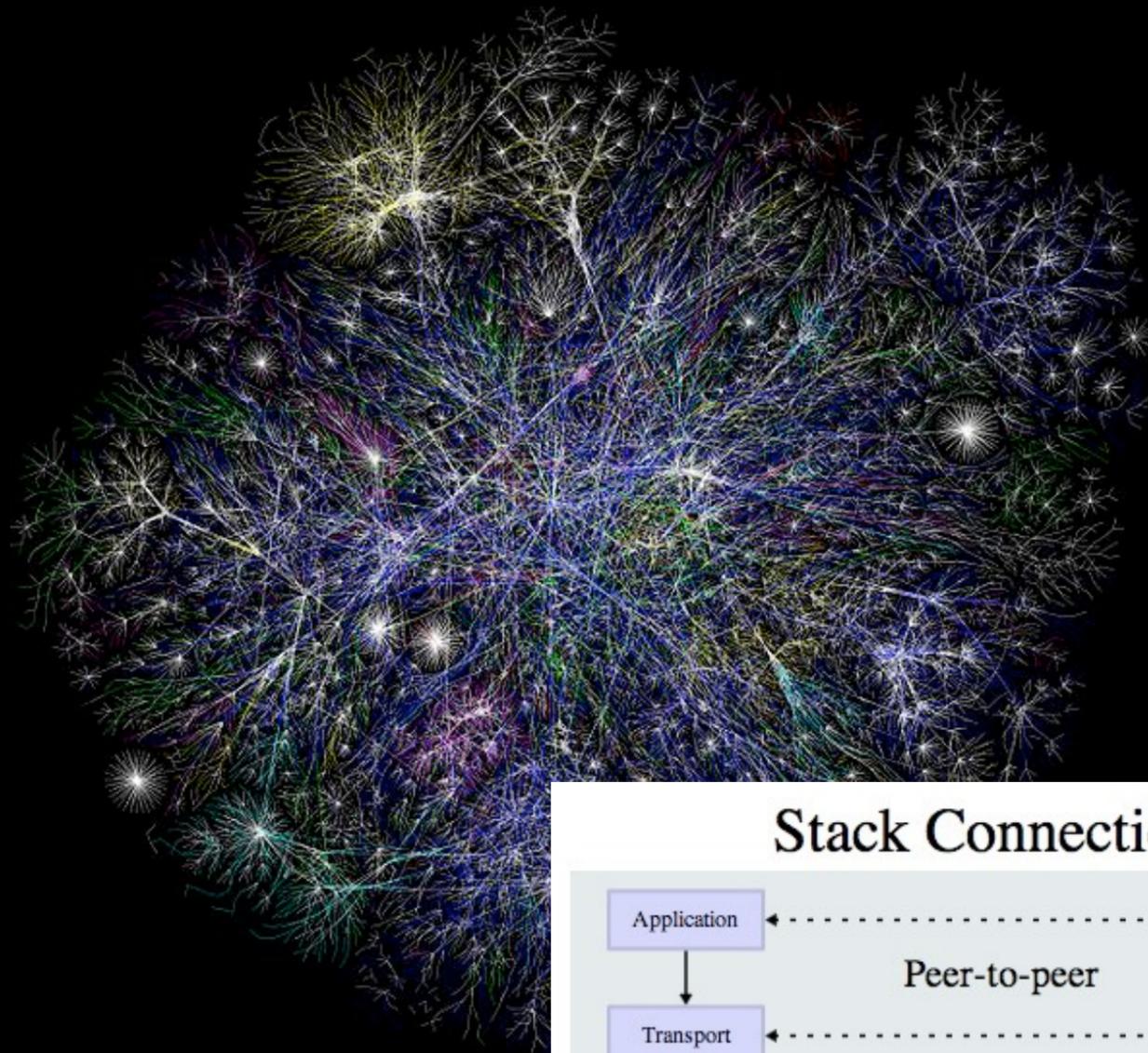
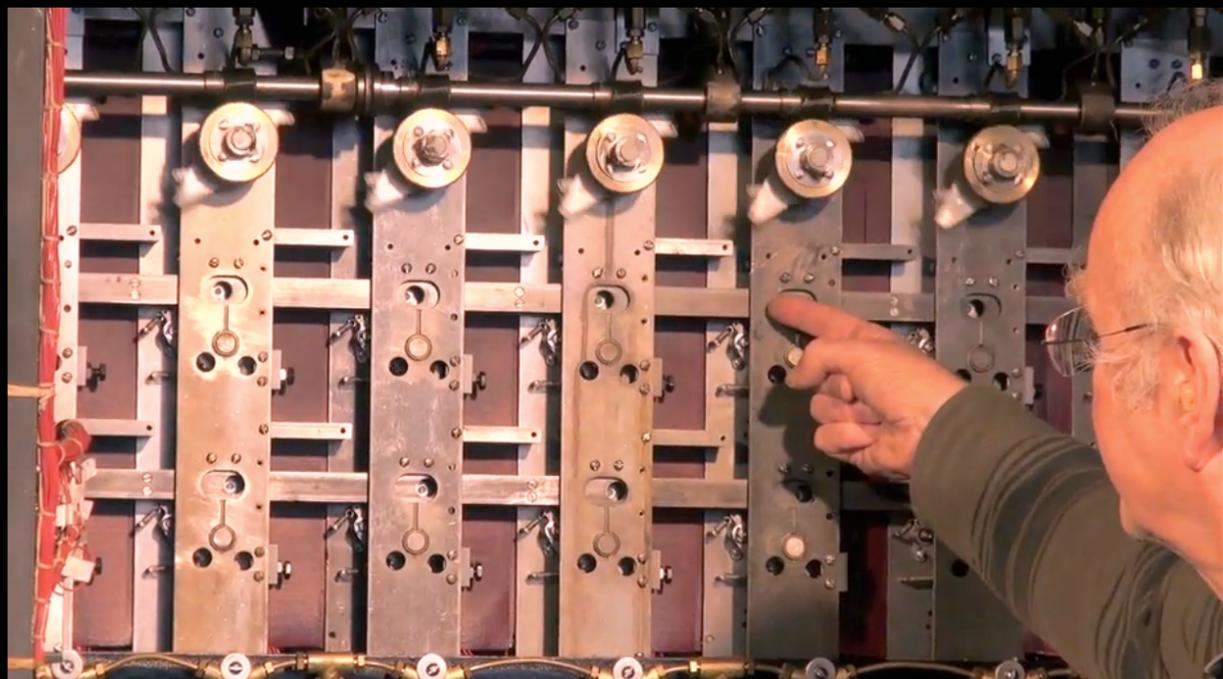
Link Layer (Ethernet, WiFi)  
Physical Connections

# The Internet: An Amazing Design

- Hundreds of millions of computers
- Thousands of routers inside the Internet
- Hundreds of millions of simultaneous connections
- Trillions of bytes of data moved per second around the world
- And it works

# The Internet

- It is said that “The Internet is the largest single engineering effort ever created by mankind”
- It was created to work in an organic way - to repair itself and automatically adjust when parts fail
- No one part of the Internet knows all of the Internet (like life)
- It is never 100% up - but it seems up all the time



# We are not done experimenting...

- There is still very active exploration on how network technology can be improved
- Content-Centric Networking is only one advanced idea
- Routers in the future can have \*lots\* of memory - lets try not to send the same piece of data more than once





# Additional Source Information

- xkcd, <http://xkcd.com/742/>, CC: BY-NC, <http://creativecommons.org/licenses/by-nc/2.5/>
- Internet Protocol Suite Diagrams: Kbrose, Wikimedia Commons, [http://upload.wikimedia.org/wikipedia/commons/c/c4/IP\\_stack\\_connections.svg](http://upload.wikimedia.org/wikipedia/commons/c/c4/IP_stack_connections.svg), CC:BY-SA, <http://creativecommons.org/licenses/by-sa/3.0/deed.en>
- All your bases are belong to me: Karin Dalziel, Flickr, <http://www.flickr.com/photos/nirak/270213335/>, CC:BY, <http://creativecommons.org/licenses/by/2.0/deed.en>
- Internet Map: The Opte Project, Wikimedia Commons, [http://upload.wikimedia.org/wikipedia/commons/d/d2/Internet\\_map\\_1024.jpg](http://upload.wikimedia.org/wikipedia/commons/d/d2/Internet_map_1024.jpg), CC:BY, <http://creativecommons.org/licenses/by/2.5/deed.en>

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